Administrivia

• Project 2 due Thursday
  - As before, free extension to midnight if you come to class
• Use the newsgroup!
  - Lots of similar questions to cs140-staff
  - TAs are too polite to tell people to ask on newsgroup
  - But everyone benefits if these questions are answered in public
  - Plus Ben Pfaff answers questions on newsgroup
• Midterm one week from today
  - Open book, open notes (but not open notebook computer)
• Review section for midterm this Friday
• Section for Project 3 next Friday

Paging

• Use disk to simulate larger virtual than physical mem

Working set model

• Disk much, much slower than memory
  - Goal: Run at memory, not disk speeds
• 90/10 rule: 10% of memory gets 90% of memory refs
  - So, keep that 10% in real memory, the other 90% on disk
  - How to pick which 10%?

Re-starting instructions

• Hardware provides kernel w. info about page fault
  - Faulting virtual address (In %cr2 reg on x86—may have seen it if you modified Pintos page fault and used fault_addr)
  - Address of instruction that caused fault
  - Was the access a read or write? Was it an instruction fetch? Was it caused by user access to kernel-only memory?
• Hardware must allow resuming after a fault
• Idempotent instructions are easy
  - E.g., simple load or store instruction can be restarted
  - Just re-execute any instruction that only accesses one address
• Complex instructions must be re-started, too
  - E.g., x86 move string instructions
  - Specify src, dst, count in %esi, %edi, %ecx registers
  - On fault, registers adjusted to resume where move left off

What to fetch

• Bring in page that caused page fault
• Pre-fetch surrounding pages?
  - Reading two disk blocks approximately as fast as reading one
  - As long as no track/head switch, seek time dominates
  - If application exhibits spacial locality, then big win to store and read multiple contiguous pages
• Also pre-zero unused pages in idle loop
  - Need 0-filled pages for stack, heap, anonymously mmapped memory
  - Zeroing them only on demand is slower
  - So many OSes zero freed pages while CPU is idle
Selecting physical pages

- May need to eject some pages
  - More on eviction policy in two slides
- May also have a choice of physical pages
- Direct-mapped physical caches
  - Virtual → Physical mapping can affect performance
  - Applications can conflict with each other or themselves
  - Scientific applications benefit if consecutive virtual pages to not conflict in the cache
  - Many other applications do better with random mapping

Superpages

- How should OS make use of “large” mappings
  - x86 has 2/4MB pages that might be useful
  - Alpha has even more choices: 8KB, 64KB, 512KB, 4MB
- Sometimes more pages in L2 cache than TLB entries
  - Don’t want costly TLB misses going to main memory
- Or have two-level TLBs
  - Want to maximize hit rate in faster L1 TLB
- OS can transparently support superpages [Navarro]
  - “Reserve” appropriate physical pages if possible
  - Promote contiguous pages to superpages
  - Does complicate evicting (esp. dirty pages) – demote

Straw man: FIFO eviction

- Evict oldest fetched page in system
- Example—reference string 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5
- 3 physical pages: 9 page faults

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Belady’s Anomaly

- More phys. mem. doesn’t always mean fewer faults

Optimal page replacement

- What is optimal (if you knew the future)?
Optimal page replacement

- What is optimal (if you knew the future)?
  - Replace page that will not be used for longest period of time
- Example—reference string 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5
- With 4 physical pages:

```
1 4
2
3
4 5
```

LRU page replacement

- Approximate optimal with least recently used
  - Because past often predicts the future
- Example—reference string 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5
- With 4 physical pages: 8 page faults

```
1 5
2
3 5 4
4 3
```

- Problem 1: Can be pessimal – example?
- Problem 2: How to implement?

Straw man LRU implementations

- Stamp PTEs with timer value
  - E.g., CPU has cycle counter
  - Automatically writes value to PTE on each page access
  - Scan page table to find oldest counter value = LRU page
  - Problem: Would double memory traffic!
- Keep doubly-linked list of pages
  - On access remove page, place at tail of list
  - Problem: again, very expensive
- What to do?
  - Just approximate LRU, don’t try to do it exactly

Clock algorithm

- Use accessed bit supported by most hardware
  - E.g., Pentium will write 1 to A bit in PTE on first access
  - Software managed TLBs like MIPS can do the same
- Do FIFO but skip accessed pages
- Keep pages in circular FIFO list
- Scan:
  - page’s A bit = 1, set to 0 & skip
  - else if A == 0, evict
- A.k.a. second-chance replacement

Clock alg. (continued)

- Large memory may be a problem
  - Most pages reference in long interval
- Add a second clock hand
  - Two hands move in lockstep
  - Leading hand clears A bits
  - Trailing hand evicts pages with A=0
- Can also take advantage of hardware Dirty bit
  - Each page can be (Unaccessed, Clean), (Unaccessed, Dirty), (Accessed, Clean), or (Accessed, Dirty)
  - Consider clean pages for eviction before dirty
- Or use n-bit accessed \textit{count} instead just A bit
  - On sweep: \textit{count} = (A \ll (n – 1)) \land (\text{count} \gg 1)
  - Evict page with lowest \textit{count}
Other replacement algorithms

- **Random eviction**
  - Dirt simple to implement
  - Not overly horrible (avoids Belady & pathological cases)

- **LFU (least frequently used) eviction**
  - Instead of just a bit, count # times each page accessed
  - Least frequently accessed must not be very useful
  - (or maybe was just brought in and is about to be used)
  - Decay usage counts over time (for pages that fall out of usage)

- **MFU (most frequently used) algorithm**
  - Because page with the smallest count was probably just brought in and has yet to be used

- Neither LFU nor MFU used very commonly

Page buffering

- **Idea:** reduce # of I/Os on the critical path
- **Keep pool of free page frames**
  - On fault, still select victim page to evict
  - But read fetched page into already free page
  - Can resume execution while writing out victim page
  - Then add victim page to free pool
- **Can also yank pages back from free pool**
  - Contains only clean pages, but may still have data
  - If page fault on page still in free pool, recycle

Page allocation

- Allocation can be **global or local**
- **Global allocation doesn’t consider page ownership**
  - E.g., with LRU, evict least recently used page of any proc
  - Works well if P1 needs 20% of memory and P2 needs 70%:
    - Doesn’t protect you from memory pigs
      (imagine P2 keeps looping through array that is size of mem)
- **Local allocation isolates processes (or users)**
  - Separately determine how much mem each proc. should have
  - Then use LRU/clock/etc. to determine which pages to evict within each process

Thrashing

- **Thrashing:** processes on system require more memory than it has
  - Each time one page is brought in, another page, whose contents will soon be referenced, is thrown out
  - Processes will spend all of their time blocked, waiting for pages to be fetched from disk
  - I/O devs at 100% utilization but system not getting much useful work done

- What we wanted: virtual memory the size of disk with access time the speed of physical memory
- What we have: memory with access time of disk

Reasons for thrashing

- Process doesn’t reuse memory, so caching doesn’t work (past != future)
- Process does reuse memory, but it does not “fit”
- Individually, all processes fit and reuse memory, but too many for system
  - At least this case is possible to address
Multiprogramming & Thrashing

- Need to shed load when thrashing

Dealing with thrashing

- Approach 1: working set
  - Thrashing viewed from a caching perspective: given locality of reference, how big a cache does the process need?
  - Or: how much memory does process need in order to make reasonable progress (its working set)?
  - Only run processes whose memory requirements can be satisfied

- Approach 2: page fault frequency
  - Thrashing viewed as poor ratio of fetch to work
  - PFF = page faults / instructions executed
  - If PFF rises above threshold, process needs more memory
  - not enough memory on the system? Swap out.
  - If PFF sinks below threshold, memory can be taken away

Working sets

- Working set changes across phases
  - Balloons during transition

Calculating the working set

- Working set: all pages proc. will access in next $T$ time
  - Can’t calculate without predicting future
- Approximate by assuming past predicts future
  - So working set $\approx$ pages accessed in last $T$ time
- Keep idle time for each page
- Periodically scan all resident pages in system
  - A bit set? Clear it and clear the page’s idle time
  - A bit clear? Add CPU consumed since last scan to idle time
  - Working set is pages with idle time $< T$

Two-level scheduler

- Divide processes into active & inactive
  - Active – means working set resident in memory
  - Inactive – working set intentionally not loaded
- Balance set: union of all active working sets
  - Must keep balance set smaller than physical memory
- Use long-term scheduler [recall from lecture 4]
  - Moves proc active $\rightarrow$ inactive until balance set small enough
  - Periodically allows inactive to become active
  - As working set changes, must update balance set
- Complications
  - How to chose idle time threshold $T$?
  - How to pick processes for active set
  - How to count shared memory (e.g., libc.so)

Some complications of paging

- What happens to available memory?
  - Some physical memory tied up by kernel VM structures
- What happens to user/kernel crossings?
  - More crossings into kernel
  - Pointers in syscall arguments must be checked
    (can’t just kill proc. if page not present—might need to page in)
- What happens to IPC?
  - Must change hardware address space
  - Increases TLB misses
  - Context switch flushes TLB entirely on old x86 machines
    (But not on MIPS... Why?)
64-bit address spaces

- Recall x86-64 only has 48-bit virtual address space
- What if you want a 64-bit virtual address space?
  - Straight hierarchical page tables not efficient
- **Solution 1: Guarded page tables [Liedtke]**
  - Omit intermediary tables with only one entry
  - Add predicate in high level tables, stating the only virtual address range mapped underneath + # bits to skip
- **Solution 2: Hashed page tables**
  - Store Virtual → Physical translations in hash table
  - Table size proportional to physical memory
  - Clustering makes this more efficient [Talluri]

Recall typical virtual address space

- Dynamically allocated memory goes in heap
- Top of heap called **breakpoint**
  - Addresses between breakpoint and stack all invalid

Early VM system calls

- OS keeps “Breakpoint” – top of heap
  - Memory regions between breakpoint & stack fault on access
- char *brk (const char *addr);
  - Set and return new value of breakpoint
- char *sbrk (int incr);
  - Increment value of the breakpoint & return old value
- Can implement `malloc` in terms of `sbrk`
  - But hard to “give back” physical memory to system

Memory mapped files

- Other memory objects between heap and stack

mmap system call

- void *mmap (void *addr, size_t len, int prot, int flags, int fd, off_t offset);
  - Map file specified by fd at virtual address addr
  - If addr is NULL, let kernel choose the address
- prot – protection of region
  - OR of PROT_EXEC, PROT_READ, PROT_WRITE, PROT_NONE
- flags
  - MAP_ANON – anonymous memory (fd should be -1)
  - MAP_PRIVATE – modifications are private
  - MAP_SHARED – modifications seen by everyone

More VM system calls

- int msync(void *addr, size_t len, int flags);
  - Flush changes of mmapped file to backing store
- int munmap(void *addr, size_t len)
  - Removes memory-mapped object
- int mprotect(void *addr, size_t len, int prot)
  - Changes protection on pages to or of PROT_...
- int mincore(void *addr, size_t len, char *vec)
  - Returns in vec which pages present
Exposing page faults

```c
typedef struct sigaction {
    union {
        void (*sa_handler)(int);
        void (*sa_sigaction)(int, siginfo_t *, void *);
    }
    sigset_t sa_mask; /* signal mask to apply */
    int sa_flags;
};

int sigaction (int sig, const struct sigaction *act, 
               struct sigaction *oact)
```

- Can specify function to run on SIGSEGV (Unix signal raised on invalid memory access)

VM tricks at user level

- **Combination of `mprotect/sigaction** very powerful
  - Can use OS VM tricks in user-level programs [Appel]
  - E.g., fault, unprotect page, return from signal handler
- **Technique used in object-oriented databases**
  - Bring in objects on demand
  - Keep track of which objects may be dirty
  - Manage memory as a cache for much larger object DB
- **Other interesting applications**
  - Useful for some garbage collection algorithms
  - Snapshot processes (copy on write)

4.4 BSD VM system [McKusick]

- Each process has a `vmspace` structure containing
  - `vm_map` – machine-independent virtual address space
  - `vm_pmap` – machine-dependent data structures
  - statistics – e.g. for syscalls like `getrusage()`
- `vm_map` is a linked list of `vm_map_entry` structs
  - `vm_map_entry` covers contiguous virtual memory
  - points to `vm_object` struct
- `vm_object` is source of data
  - e.g. vnode object for memory mapped file
  - points to list of `vm_page` structs (one per mapped page)
  - `shadow objects` point to other objects for copy on write

Pmap (machine-dependent) layer

- Pmap layer holds architecture-specific VM code
- VM layer invokes pmap layer
  - On page faults to install mappings
  - To protect or unmap pages
  - To ask for dirty/accessed bits
- Pmap layer is lazy and can discard mappings
  - No need to notify VM layer
  - Process will fault and VM layer must reinstall mapping
- Pmap handles restrictions imposed by cache
Example uses

- `vm_map_entry` structs for a process
  - r/o text segment → file object
  - r/w data segment → shadow object → file object
  - r/w stack → anonymous object

- New `vm_map_entry` objects after a fork:
  - Share text segment directly (read-only)
  - Share data through two new shadow objects
    (must share pre-fork but not post fork changes)
  - Share stack through two new shadow objects

- Must discard/collapse superfluous shadows
  - E.g., when child process exits

What happens on a fault?

- Traverse `vm_map_entry` list to get appropriate entry
  - No entry? Protection violation? Send process a SIGSEGV

- Traverse list of [shadow] objects

- For each object, traverse `vm_page` structs

- Found a `vm_page` for this object?
  - If first `vm_object` in chain, map page
  - If read fault, install page read only
  - Else if write fault, install copy of page

- Else get page from object
  - Page in from file, zero-fill new page, etc.