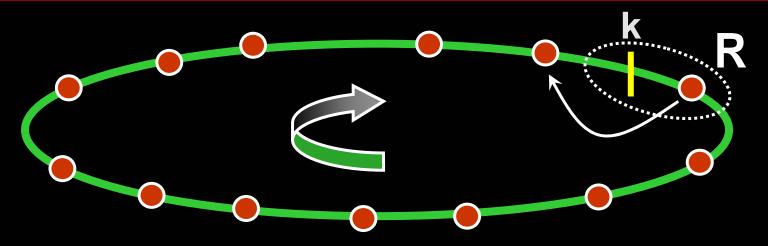


# Non-Transitive Connectivity and DHTs

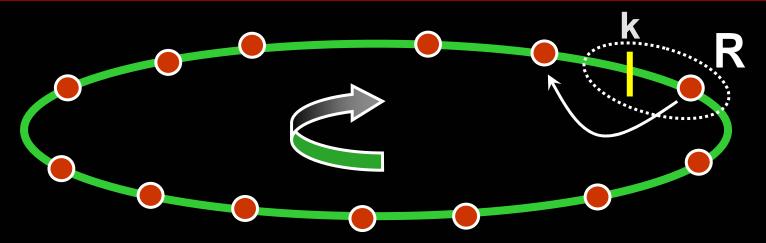
### Mike Freedman Karthik Lakshminarayanan Sean Rhea Ion Stoica

WORLDS 2005



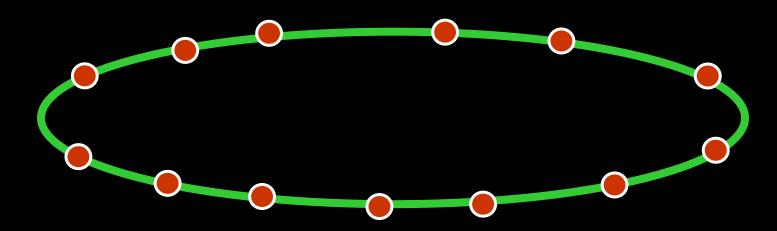
System assigns keys to nodes
All nodes agree on assignment

Chord assigns keys as integers modulo 2<sup>160</sup>
Assigns keys via successor relationship
Each node must know predecessor

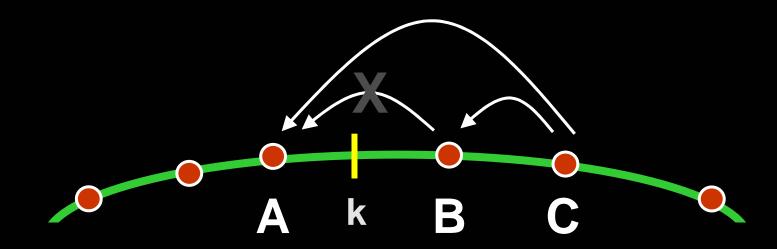


Used to store and retrieve (key, value) pairs

- Any node can discover key's successor, yet without full knowledge of network
  - Implies some form of routing



### All have implicit assumption: full connectivity



- All have implicit assumption: full connectivity
- Non-transitive connectivity (NTC) not uncommon

 $B \leftrightarrow C$  ,  $C \leftrightarrow A$  ,  $A \not\leftrightarrow B$ 

A thinks C is its successor!

## Does non-transitivity exist?

- Gerding/Stribling PlanetLab study
  - 9% of all node triples exhibit NTC
  - Attributed high extent to Internet-2
- Yet NTC is also transient
  - One 3 hour PlanetLab all-pair-pings trace
  - 2.9% have persistent NTC
  - 2.3% have intermittent NTC
  - 1.3% fail only for a single 15-minute snapshot

Level3 ↔ Cogent, but Level3 ↔ X ↔ Cogent
 NTC motivates RON, Detour, and SOSR!

# Our contributions

We have built and run Bamboo (OpenDHT),
 Chord (i3), Kademlia (Coral) for > 1 year

Vanilla DHT algorithms break under NTC

 Identify four main algorithmic problems and present our solutions

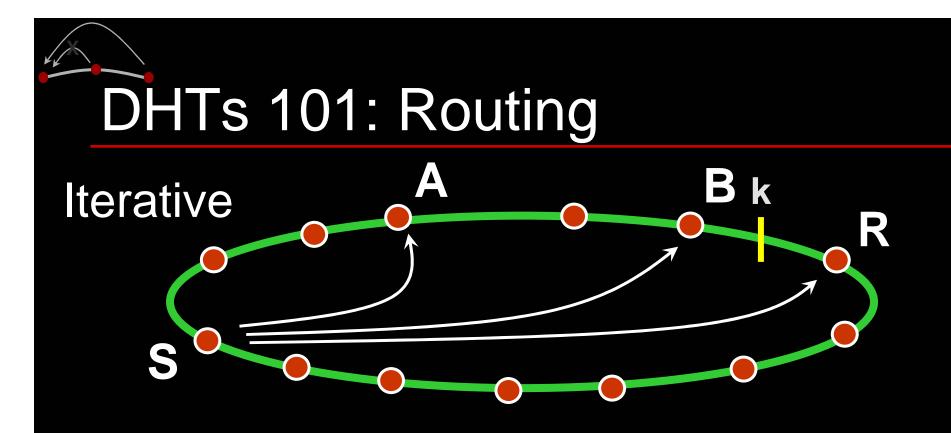
# Our goals

### Short-term

- Inform other developers about NTC solutions
- Important: DHTs are being widely deployed in Overnet, Morpheus, and BitTorrent

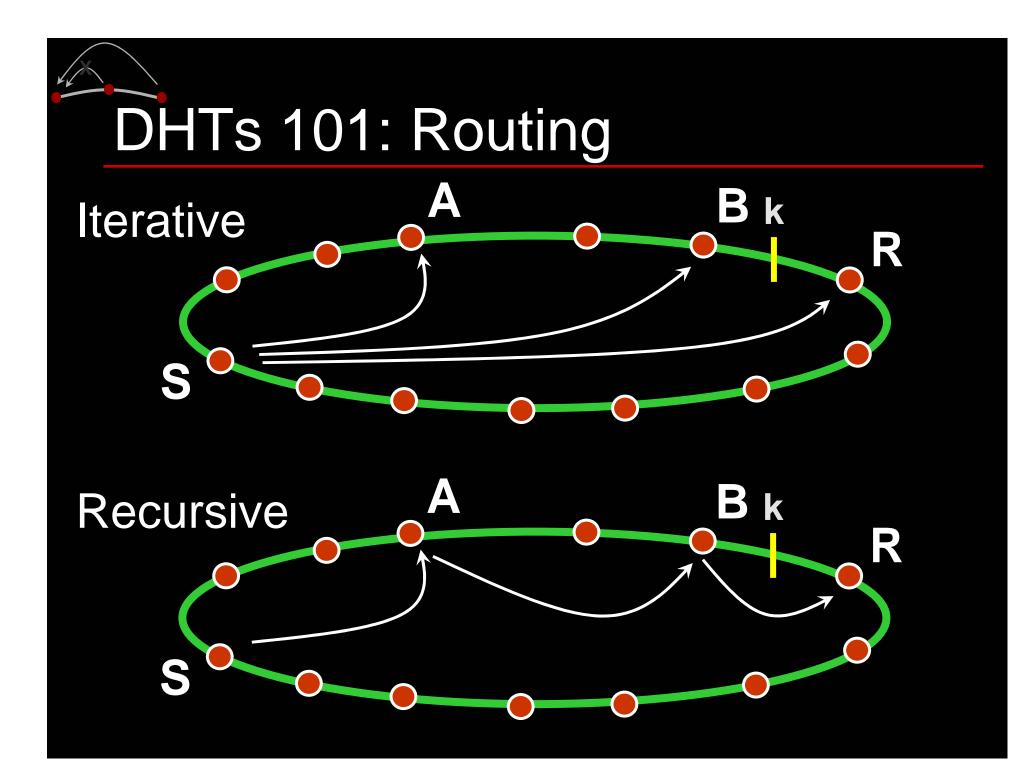
### Long-term

- Encourage new designs to directly handle NTC
- (This topic is far from solved)

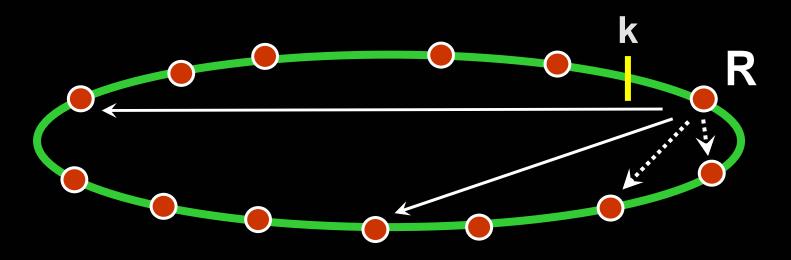


Key space defines an identifier distance

 Routing ideally proceeds by halving distance to destination per overlay hop



# DHTs 101: Routing tables



successors / leaf set: ensure correctness

fingers / routing table: efficient routing
 O (log (n)) hops, generally

# Problems we identify

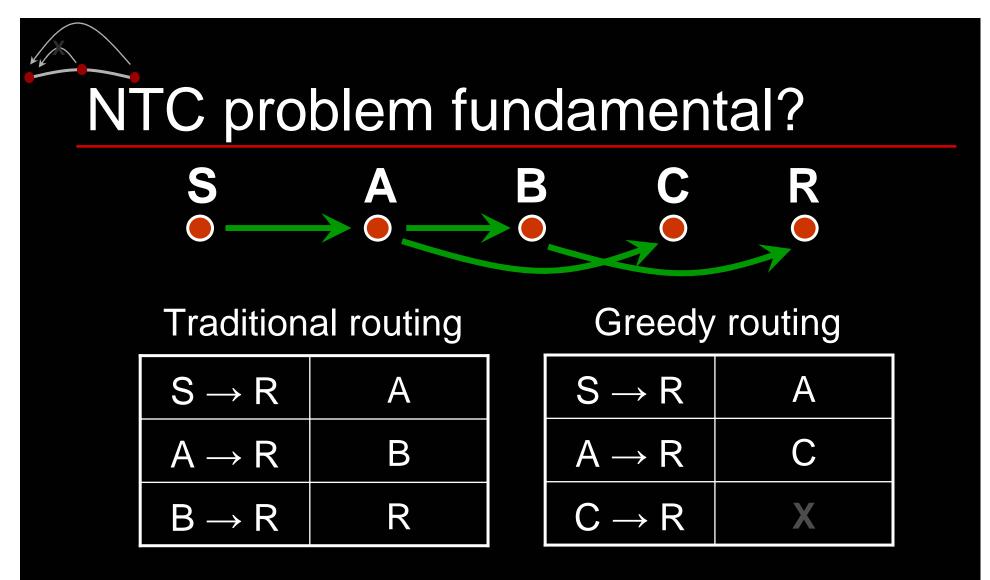
- Invisible nodes
- Routing loops
- Broken return paths
- Inconsistent roots

## NTC problem fundamental?

# S A B C R

#### **Traditional routing**

$S \rightarrow R$	A
$A \rightarrow R$	B
$B \to R$	R

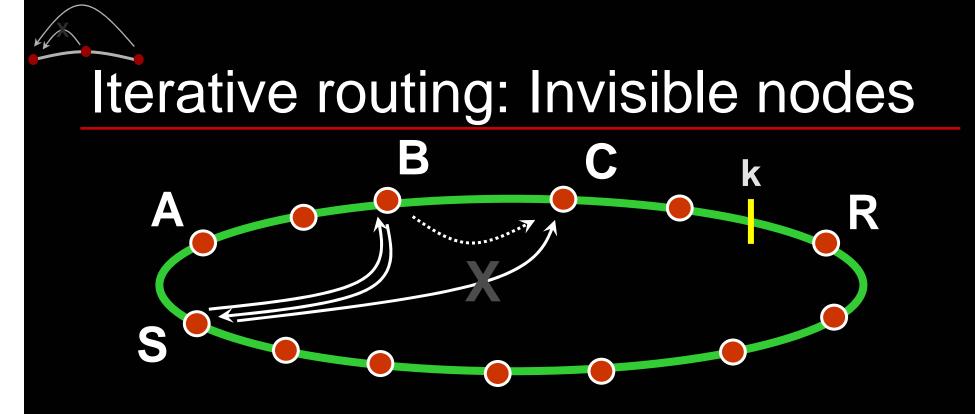


- DHTs implement greedy routing for scalability
- Sender might not use path, even though exists: finds local minima when id-distance routing

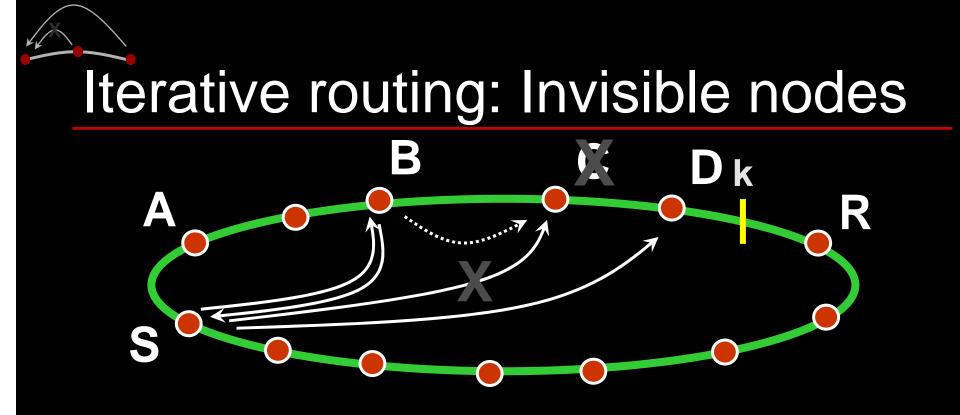
# Problems we identify

- Invisible nodes
- Routing loops
- Broken return paths
- Inconsistent roots

(First discuss how problems apply to iterative routing, then consider recursive routing.)



Invisible nodes cause lookup to halt

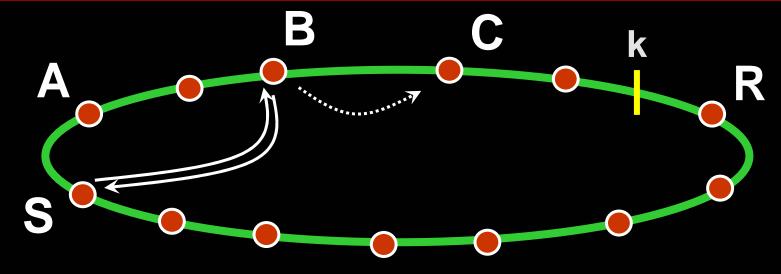


Invisible nodes cause lookup to halt

### Enable lookup to continue

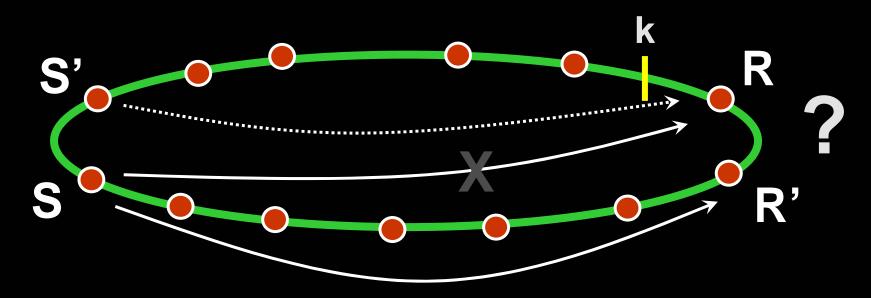
- Tighter timeouts via network coordinates
- Lookup RPCs in parallel
- Unreachable node cache

# Routing table pollution



- Many proposals for maintaining routing tables
   E.g., replace nodes with larger RTT
- Must first prevent routing table pollution
  - Only add new nodes upon contacting *directly*
  - Do not immediately remove nodes from hearsay

### Inconsistent roots



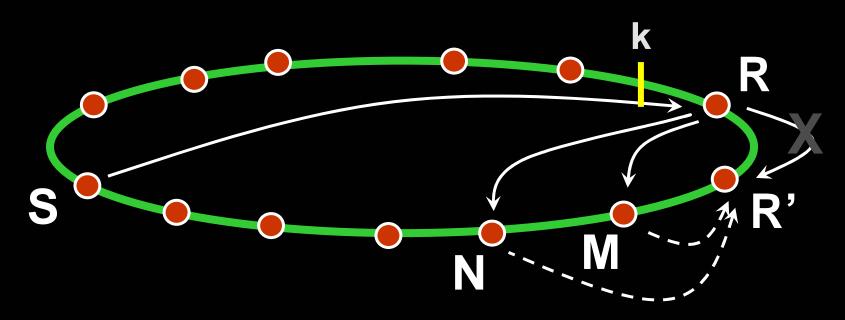
- Nodes do not agree where key is assigned: inconsistent views of root
  - Can be caused by membership changes
  - Also due to non-transitive connectivity
    - May persist indefinitely

## Inconsistent roots

- No solution when network partitions
- If non-transitivity is limited:
  - Consensus among leaf set?
    - [Etna, Rosebud]
    - Expensive in messages and bandwidth
  - Link-state routing among leaf set?
    - [Pastry 1.4.1]

Can use application-level solutions!

## Inconsistent roots



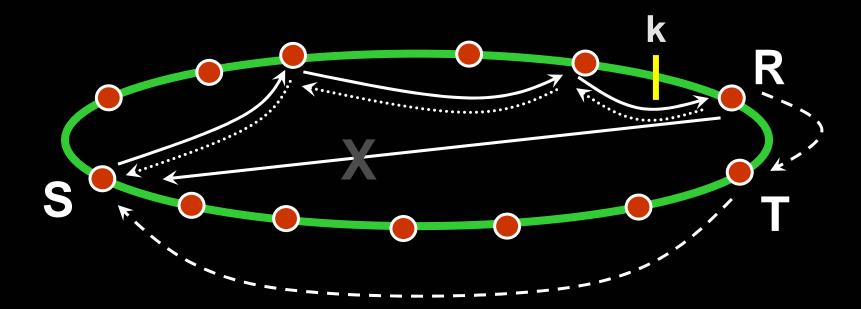
- Root replicates (key,value) among leaf set
  - Leafs periodically synchronize
  - Get gathers results from multiple leafs
  - [OpenDHT, DHash]

Not applicable when require fast update (i3)

## **Recursive routing**

- Invisible nodes
  - Must also prevent routing table pollution
  - Easier to achieve accurate timeouts
  - Harder to perform concurrent RPCs
- Inconsistent Roots
  - Similar solutions
- (Routing Loops)
- One new problem...

# Broken return paths



..........

Direct path back from R to S fails

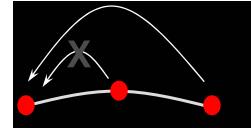
Source-route reverse path

Use single intermediate hop --

RON, Detour, SOSR...

# Summary

- Non-transitive connectivity exists
  - DHTs must deal with it
- Discovered problems the "hard way"
  - OpenDHT / Bamboo, i3 / Chord, Coral / Kademlia
  - Presented our "from the trenches" fixes
- NTC should be considered during design phase



# Thanks...

### Watch Our Real, Large Distributed Systems...

coralcdn.org opendht.org i3.cs.berkeley.edu