

The Final Nail in WEP's Coffin

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Wired Equivalent Privacy

Introduction Fragmentation Attack Implementation Conclusion



WEP is the 802.11 standard for encryption.

- Pre-shared key for whole network.
- Protects data privacy since data is encrypted.
- Access control: need key to transmit.

In practice, only half of the networks are encrypted.

• In the subset of encrypted networks, WEP is most adopted.

Popularity (%) of WEP and its alternatives based on our survey

Region	WEP	WPA	802.11i
London	76	20	4
Seattle region	85	14	1

Goals when attacking WEP

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- Decrypt data in packets.
- Obtain access to the network by being able to transmit data.
- Recover the WEP key.





Today, millions of packets are required to break a WEP key.

Our fragmentation attack allows:

- Transmitting arbitrary data after eavesdropping a single data packet on an 802.11 WEP protected network.
- 2 Real-time decryption given that network is connected to Internet.

Outline

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Introduction
 WEP Description

- WEP Attacks
- Pragmentation Attack
 Transmission
 - Decryption
- Implementation
 - Performance Evaluation

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WEP operation

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History of WEP attacks ... and how the *real* problem was ignored

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Year 2000. Keystream attacks (independent of WEP key):

- Design flaw: WEP allows keystream reuse.
- Attacks were thought impractical:
 - Need plain-text to recover keystream.
 - $\,\circ\,$ Need 2^{24} keystreams to decrypt all possible packets.

Year 2001. Weak IV attacks (recover WEP key):

- Need millions of packets. Could take hours, and usually, days.
- Use EAP-based solutions to re-key, say, every ten minutes.

Year 2006. Our contribution: fragmentation attack.

• Keystream attack which may be performed within minutes.



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Fragmentation attack

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Transmission

- Recover a keystream.
- 2 Reuse the keystream to send arbitrary data.

Keystream-based decryption

- Resend data through the AP to a buddy on the Internet.
- Recover the keystream used for encrypting the packet.

WEP key recovery

Use transmission ability for speeding up weak IV attacks.



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802.11 head	er 0×AA	0xAA	0×03	0×00	0×00	0×00	0×08	0×00			
\oplus											
802.11 head	er	Cipher-text									
	=										
		8 bytes of keystream									

Can recover 8 bytes of keystream by eavesdropping a packet.

• Can encrypt (and transmit) 8 bytes of arbitrary data.



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Decrypt data locally in linear time with respect to its length:

- If *n* bytes of keystream were known, *n* bytes of data could be decrypted. Base case: 8 bytes of keystream are known.
- Guess keystream byte *n* + 1 and verify it. Send a broadcast using extended keystream. If AP relayed it, guess is correct.
- Continue keystream expansion for whole length of packet.





To decrypt data in real-time, resend it to the Internet.

- Eavesdrop a payload to decrypt.
- ② Send two 802.11 fragments: an IP header with our buddy as destination, and the original encrypted payload as a fragment.
- 3 AP will decrypt and send it in clear-text to our Internet buddy.



Fragmentation attack Summary

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To encrypt data:

- Eavesdrop a data packet.
- ② Cipher-text \oplus known plain-text = 8 bytes of keystream.
- 3 Transmit data in multiple 8 byte fragments.

To decrypt data:

- Eavesdrop packet to decrypt.
- ② Send two 802.11 fragments:
 - In An IP header destined to a buddy on the Internet.
 - 2 A fragment containing the original eavesdropped payload.
- 3 Internet buddy will receive the payload in clear-text.

Implementation

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Hardware: Atheros chipset.

- Software radio. Ideal for packet injection.
- Supports 802.11{a,b,g}.

Software:

- FreeBSD. Added packet injection support to ath driver.
- wesside. Proof-of-concept fragmentation attack tool.

wesside

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- I Eavesdrops data packet and uses fragmentation to transmit.
- 2 Determines the network IP via keystream expansion.
- 3 Contacts buddy on Internet instructing him to flood the WiFi.
- ④ Recovers WEP key via weak IV attack (using aircrack).



Performance

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Fragmentation attack, after eavesdropping one packet:

- Recover 1500 bytes of keystream: < 2 seconds.
- Decrypt network's IP: < 30 seconds.







Fragmentation attack:

- May be performed instantly. Frequent re-keying (EAP) does not mitigate the problem. Migrate to 802.11i.
- Non-solution: ship hardware with no fragmentation support.
- Solution: ship hardware with no WEP support.

WEP history:

- Attacks evolve over time. In 2000, theoretical issues were identified. Today, we provide a practical exploit for them.
- Theoretical guidelines must be followed. Perfect example of damage incurred by keystream reuse and no authentication.





Do not use WEP-teach about its failures.

Future Work: The First Nail in WPA's Coffin...