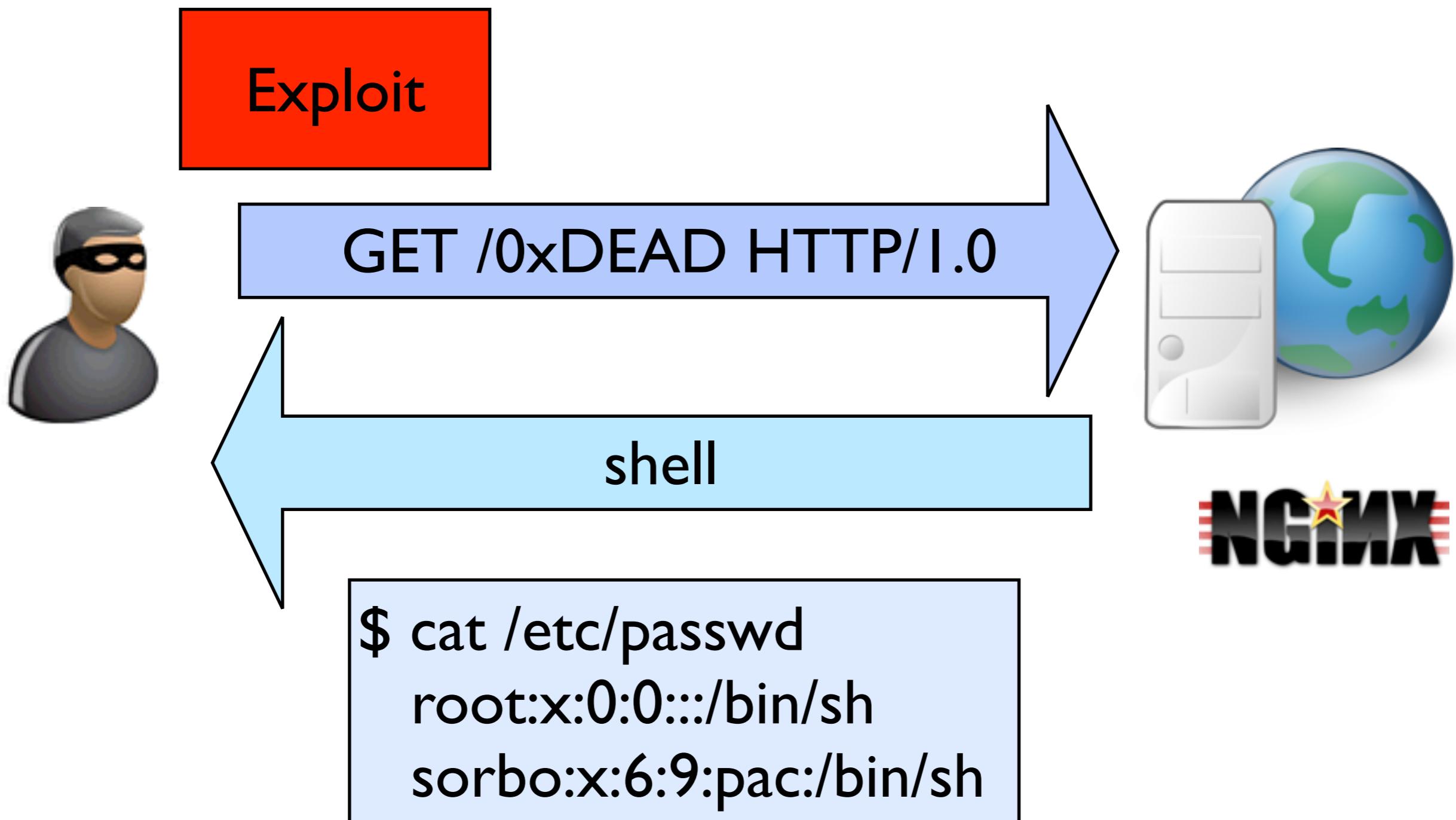


Hacking Blind

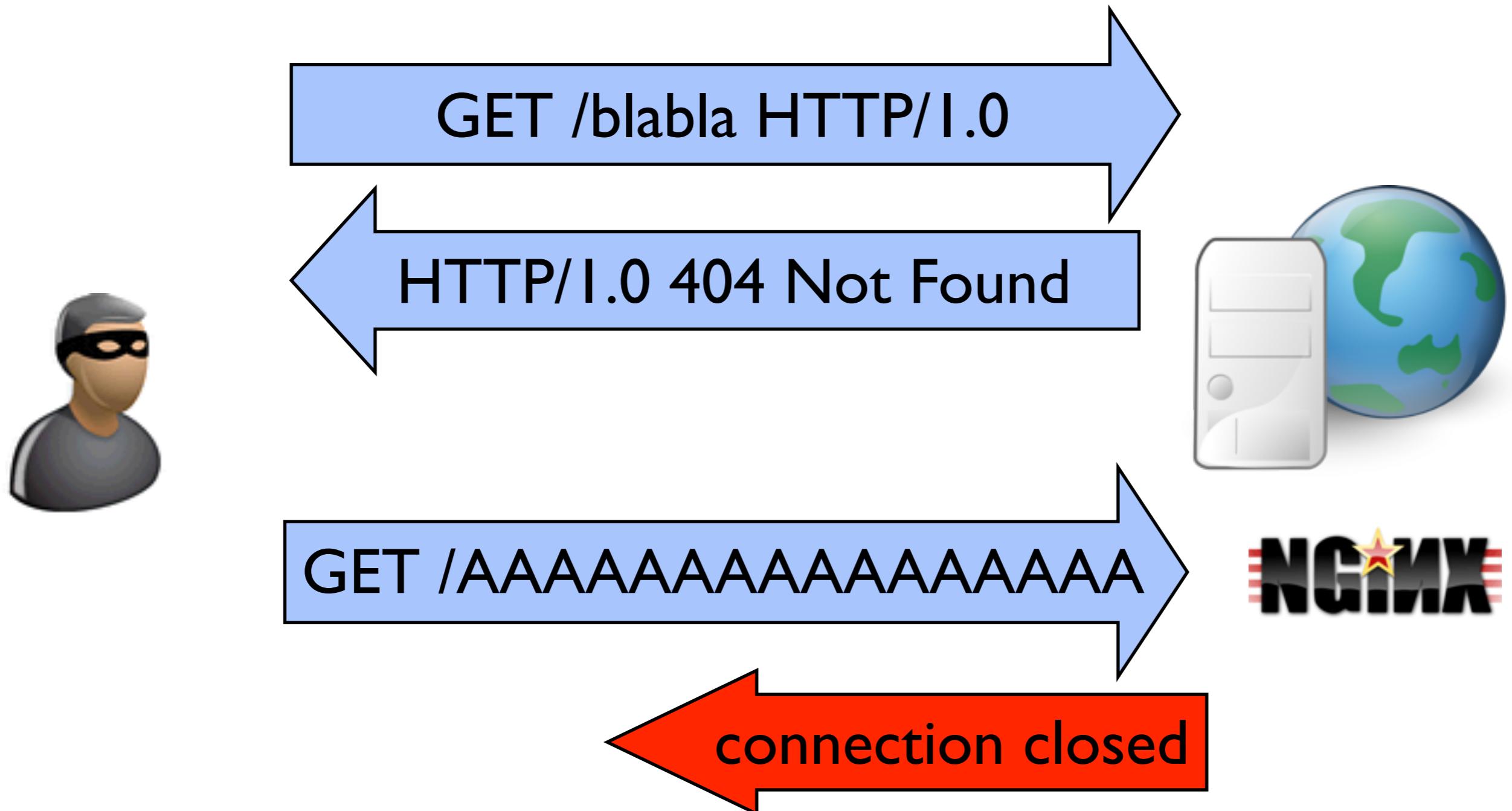
**Andrea Bittau, Adam Belay, Ali Mashtizadeh,
David Mazières, Dan Boneh**

Stanford University

Hacking buffer overflows



Crash or no Crash? Enough to build exploit



Don't even need to know what application is running!

Exploit scenarios:

1. Open source



2. Open binary



3. Closed-binary (and source)



Attack effectiveness

- Works on 64-bit Linux with ASLR, NX and canaries

Server	Requests	Time (mins)
nginx	2,401	1
MySQL	3,851	20
Toy proprietary service (unknown binary and source)	1,950	5

Attack requirements

1. Stack vulnerability, and knowledge of how to trigger it.
2. Server process that respawns after crash
 - E.g., nginx, MySQL, Apache, OpenSSH, Samba.

Outline

- Introduction.
- Background on exploits.
- Blind ROP (BROP).
- Optimizations.

Stack vulnerabilities

```
void process_packet(int s) {  
    char buf[1024];  
    int len;  
  
    read(s, &len, sizeof(len));  
    read(s, buf, len);  
  
    return;  
}
```

handle_client()

Stack:

return address
0x400000

buf[1024]

Stack vulnerabilities

```
void process_packet(int s) {  
    char buf[1024];  
    int len;  
  
    read(s, &len, sizeof(len));  
    read(s, buf, len);  
  
    return;  
}
```

handle_client()

Stack:

return address
0x400000

AAAAAA
AAAAAA
AAAAAA
AAAAAA

Stack vulnerabilities

```
void process_packet(int s) {  
    char buf[1024];  
    int len;  
  
    read(s, &len, sizeof(len));  
    read(s, buf, len);  
  
    return;  
}
```

Stack:

return address
0x41414141

AAAAAAA
AAAAAAA
AAAAAAA
AAAAAAA

??

Stack vulnerabilities

```
void process_packet(int s) {  
    char buf[1024];  
    int len;  
  
    read(s, &len, sizeof(len));  
    read(s, buf, len);  
  
    return;  
}
```

Shellcode:

```
dup2(sock, 0);  
dup2(sock, 1);  
execve("/bin/sh", 0, 0);
```

Stack:

return address
0x500000

AAAAAAA
AAAAAAA
AAAAAAA
AAAAAAA

Stack vulnerabilities

```
void process_packet(int s) {  
    char buf[1024];  
    int len;  
  
    read(s, &len, sizeof(len));  
    read(s, buf, len);  
  
    return;  
}
```

Shellcode:

```
dup2(sock, 0);  
dup2(sock, 1);  
execve("/bin/sh", 0, 0);
```

Stack:

return address
0x600000

0x1029827189
123781923719
823719287319
879181823828

Exploit protections

```
void process() {  
    char buffer[1024];  
    int len = strlen(buffer);  
  
    // ...  
  
    return; // Return address  
}  
  
return; // Shellcode:  
  
dup2(sock, 0);  
dup2(sock, 1);  
execve("/bin/sh", 0, 0);
```

Stack:

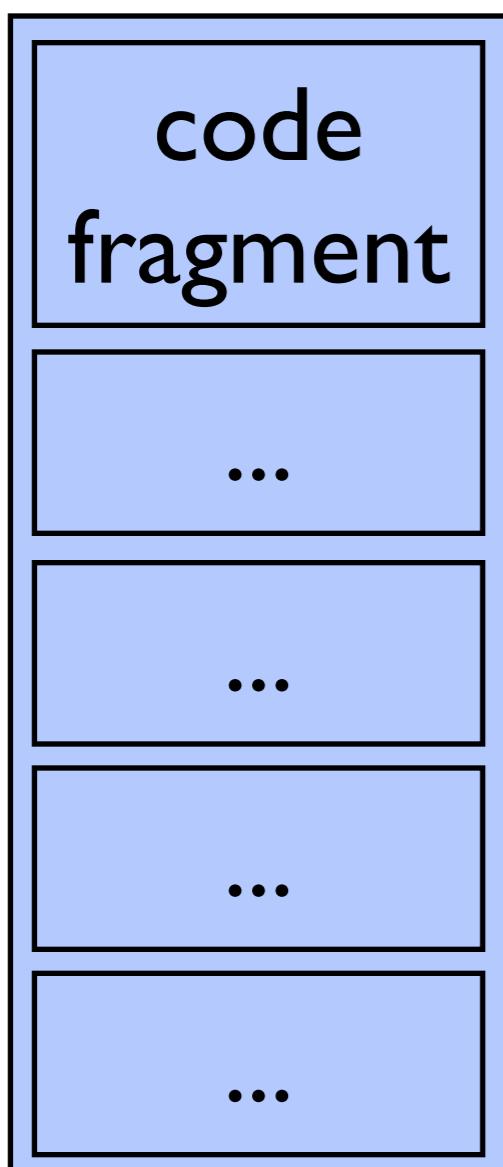
return address
0x600000

0x1029827189
123781923719
823719287319
879181823828

The diagram illustrates two exploit protection mechanisms and their bypasses. A yellow callout box labeled "1. Make stack non-executable (NX)" points to the "return" instruction, which is highlighted in red. Another yellow callout box labeled "2. Randomize memory addresses (ASLR)" points to the shellcode, also highlighted in red. Red arrows from these boxes point to a blue box labeled "Stack:" containing the return address and shellcode, illustrating how both protections are bypassed.

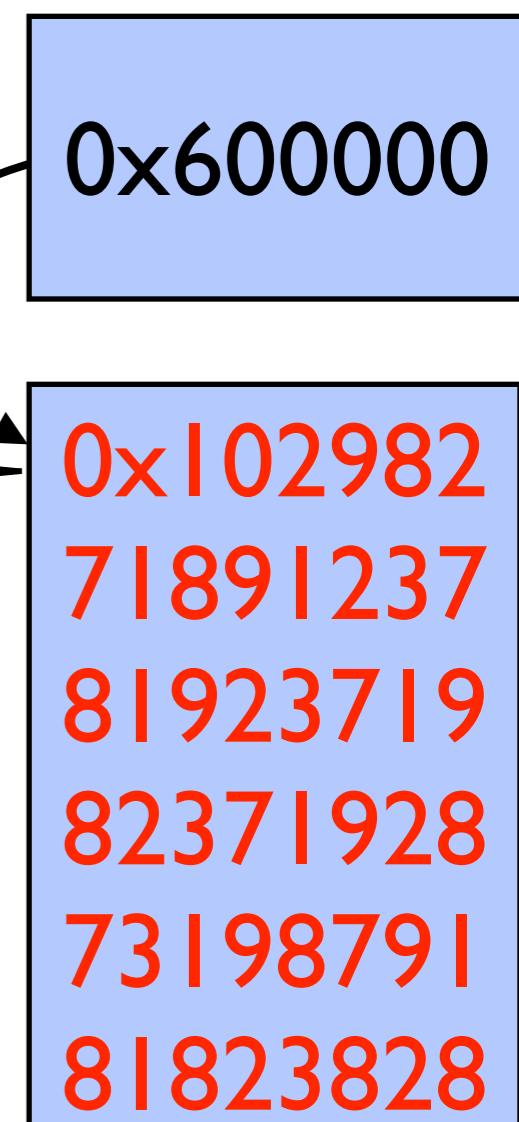
Return-Oriented Programming (ROP)

.text:



dup2(sock, 0);
dup2(sock, 1);
execve("/bin/sh", 0, 0);

Stack:

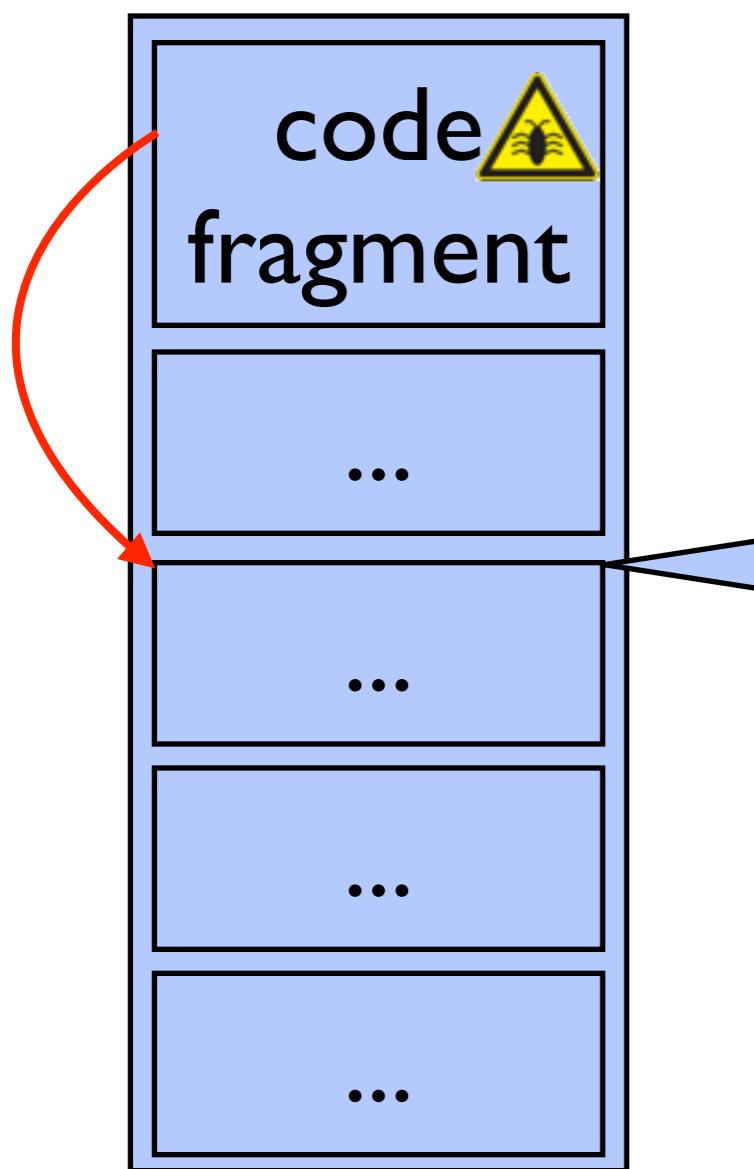


Executable

Non-Executable

Return-Oriented Programming (ROP)

.text:



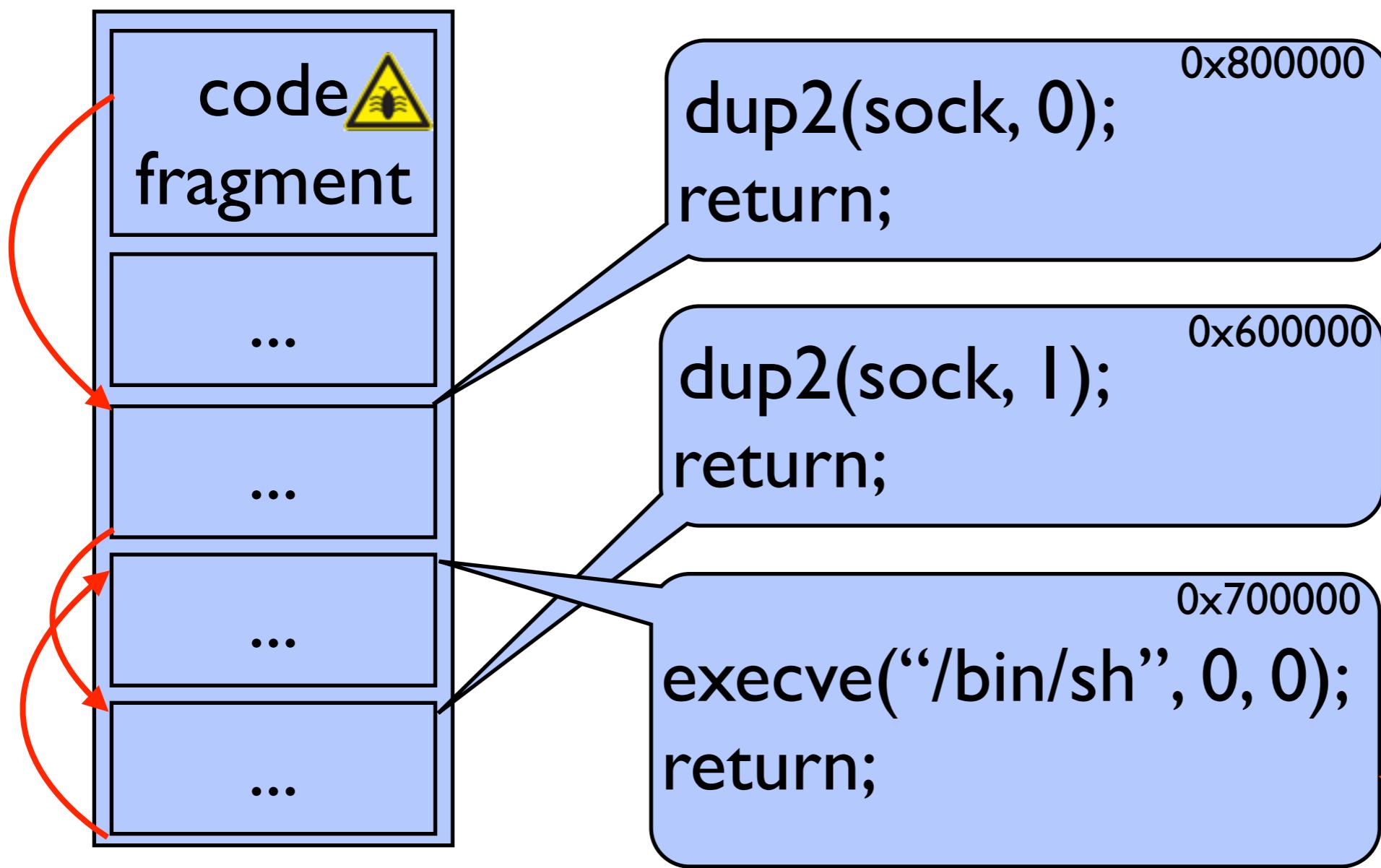
```
dup2(sock, 0);
dup2(sock, 1);
execve("/bin/sh", 0, 0);
```

Stack:

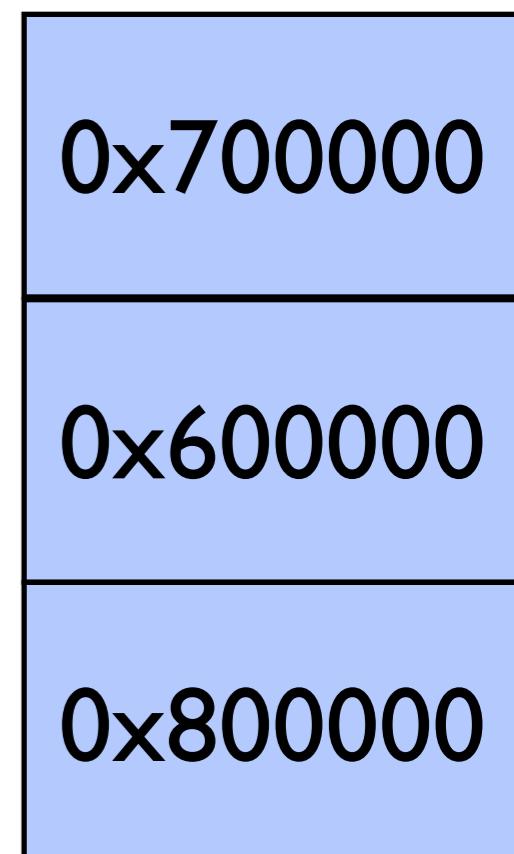


Return-Oriented Programming (ROP)

.text:



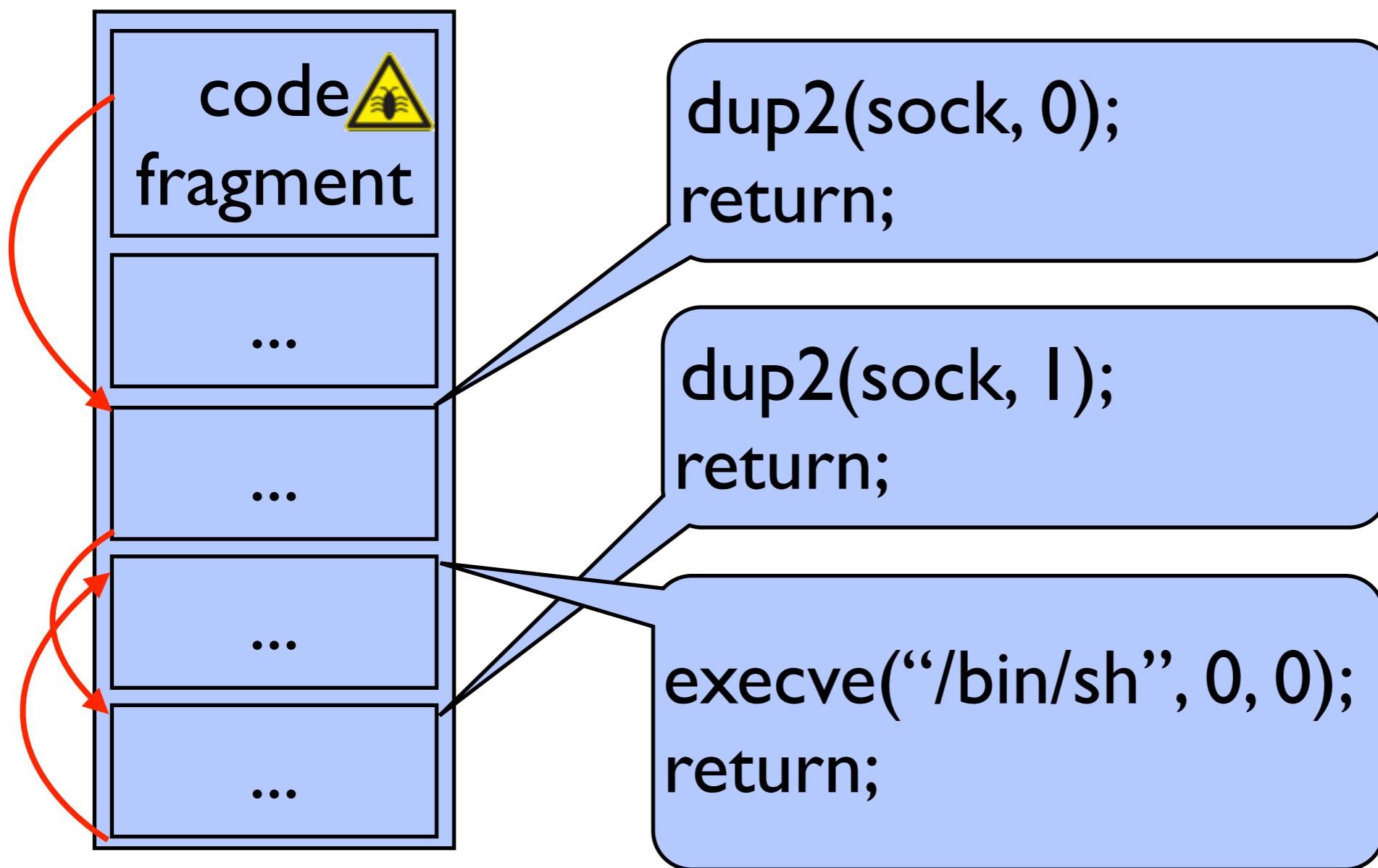
Stack:



ROP gadget

Address Space Layout Randomization (ASLR)

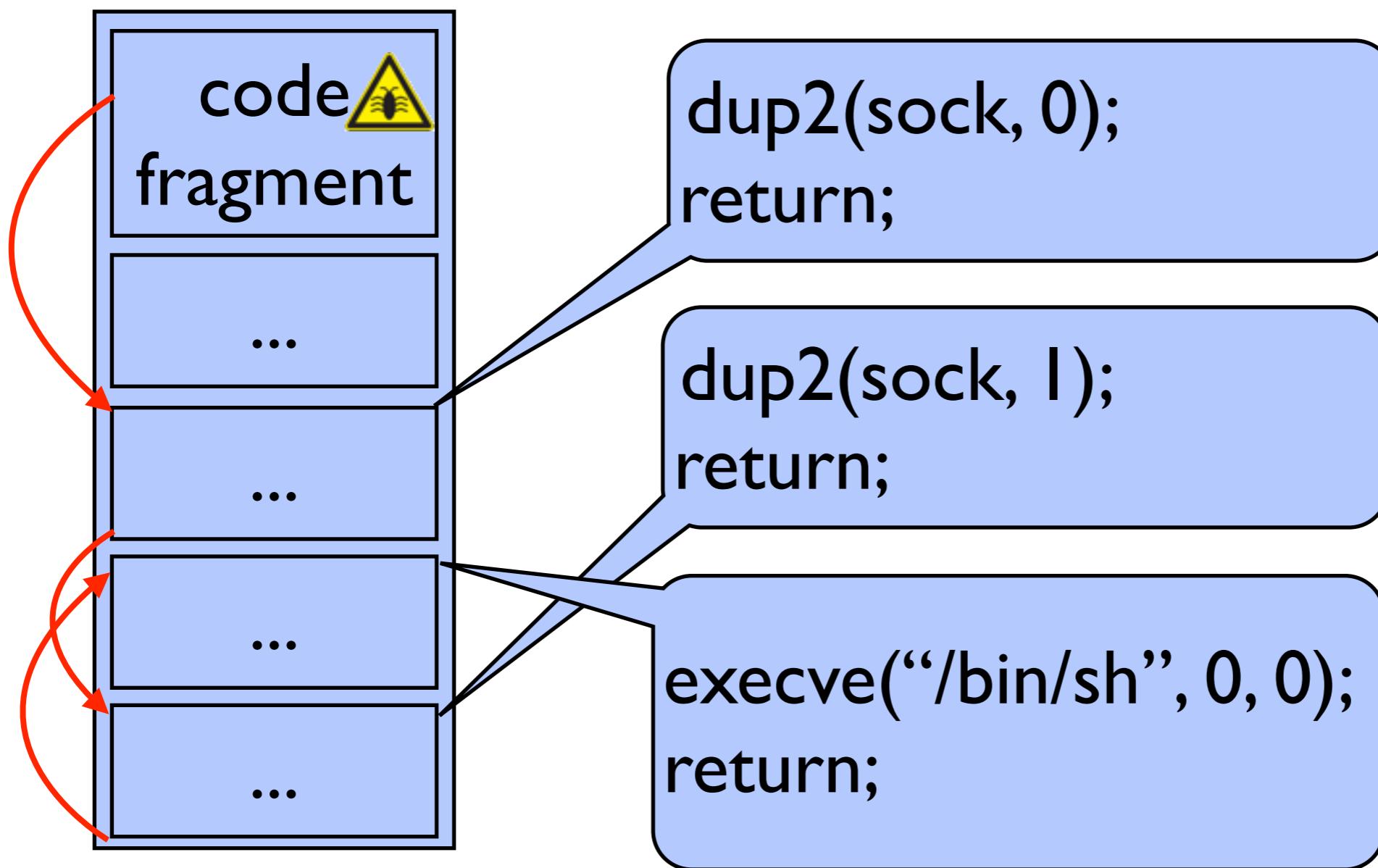
.text: 0x400000



Stack:

Address Space Layout Randomization (ASLR)

.text: 0x400000 + ??



Stack:

Exploit requirements today

1. Break ASLR.
2. Copy of binary (find ROP gadgets / break NX).
 - Is it even possible to hack unknown applications?

Blind Return-Oriented Programming (BROP)

1. Break ASLR.
2. Leak binary:
 - Remotely find enough gadgets to call `write()`.
 - `write()` binary from memory to network to disassemble and find more gadgets to finish off exploit.

Defeating ASLR: stack reading

- Overwrite a single byte with value X:
 - No crash: stack had value X.
 - Crash: guess X was incorrect.
- Known technique for leaking canaries.

Return address

buf[1024]

0x401183

Defeating ASLR: stack reading

- Overwrite a single byte with value X:
 - No crash: stack had value X.
 - Crash: guess X was incorrect.
- Known technique for leaking canaries.

Return address

00000000000000000000000000000000

0x401183

Defeating ASLR: stack reading

- Overwrite a single byte with value X:
 - No crash: stack had value X.
 - Crash: guess X was incorrect.
- Known technique for leaking canaries.

Return address

00000000000000000000000000000000

0x001183

(Was: 0x401183)

Defeating ASLR: stack reading

- Overwrite a single byte with value X:
 - No crash: stack had value X.
 - Crash: guess X was incorrect.
 - Known technique for leaking canaries.

Return address

00000000000000000000000000000000

0x01 | 183

(Was: 0x40 || 83)

Defeating ASLR: stack reading

- Overwrite a single byte with value X:
 - No crash: stack had value X.
 - Crash: guess X was incorrect.
- Known technique for leaking canaries.

Return address

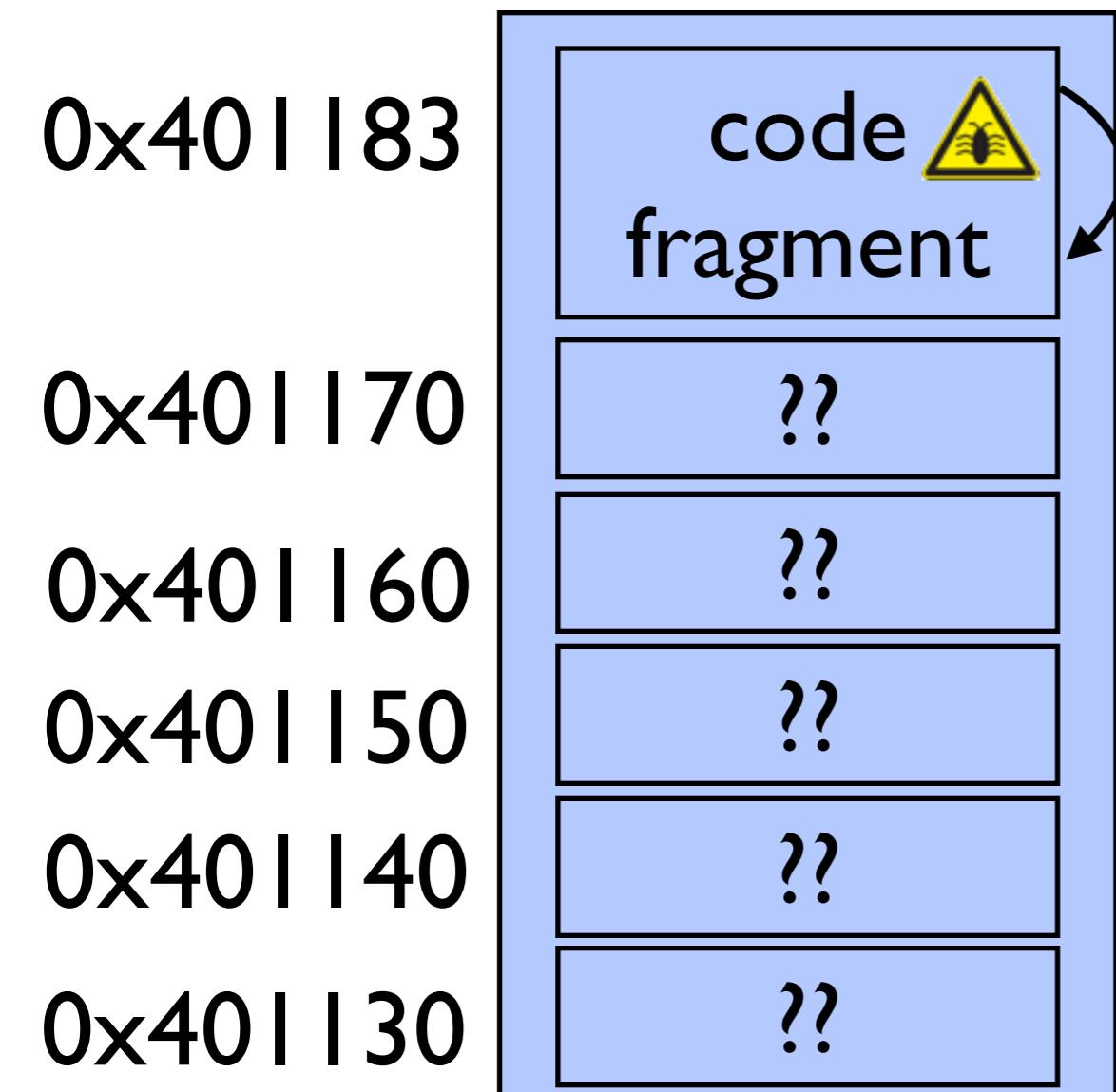
00000000000000000000000000000000

0x401183

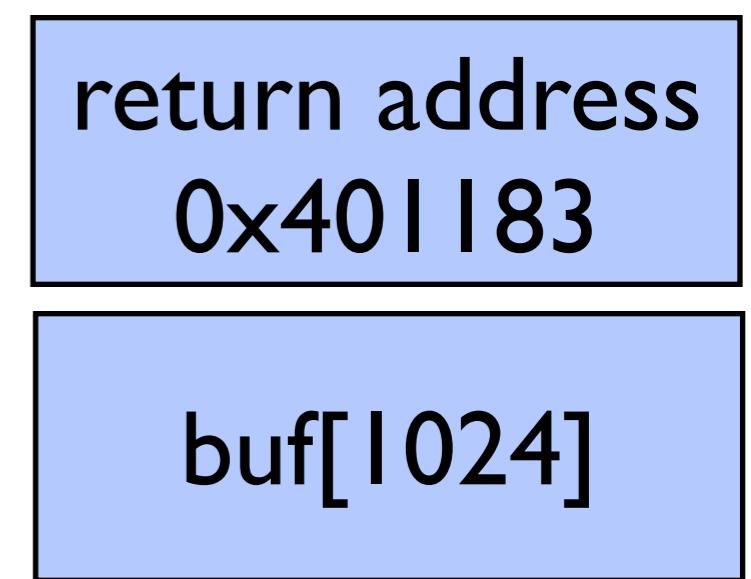
(Was: 0x401183)

How to find gadgets?

.text:

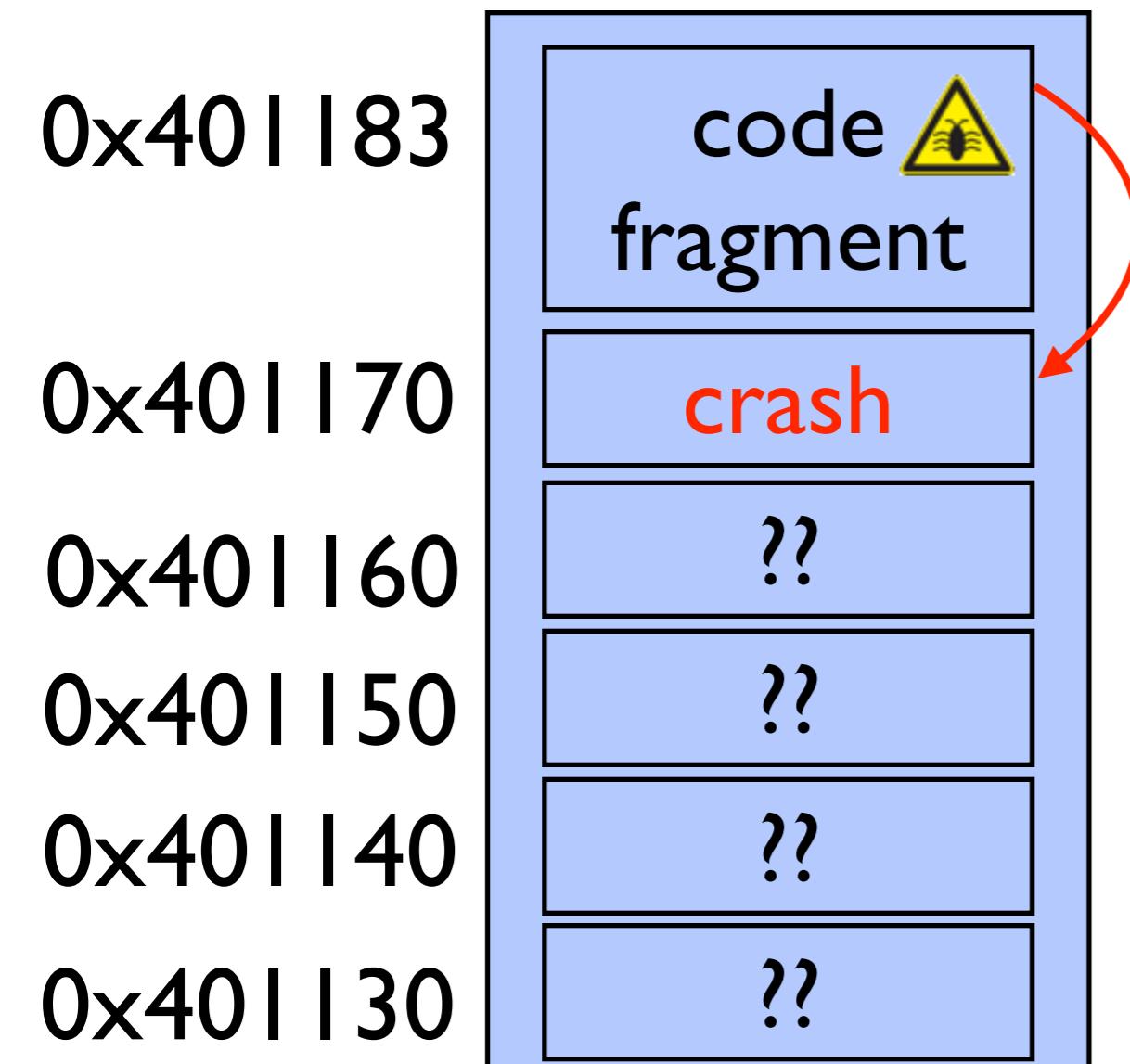


Stack:



How to find gadgets?

.text:



Connection closes

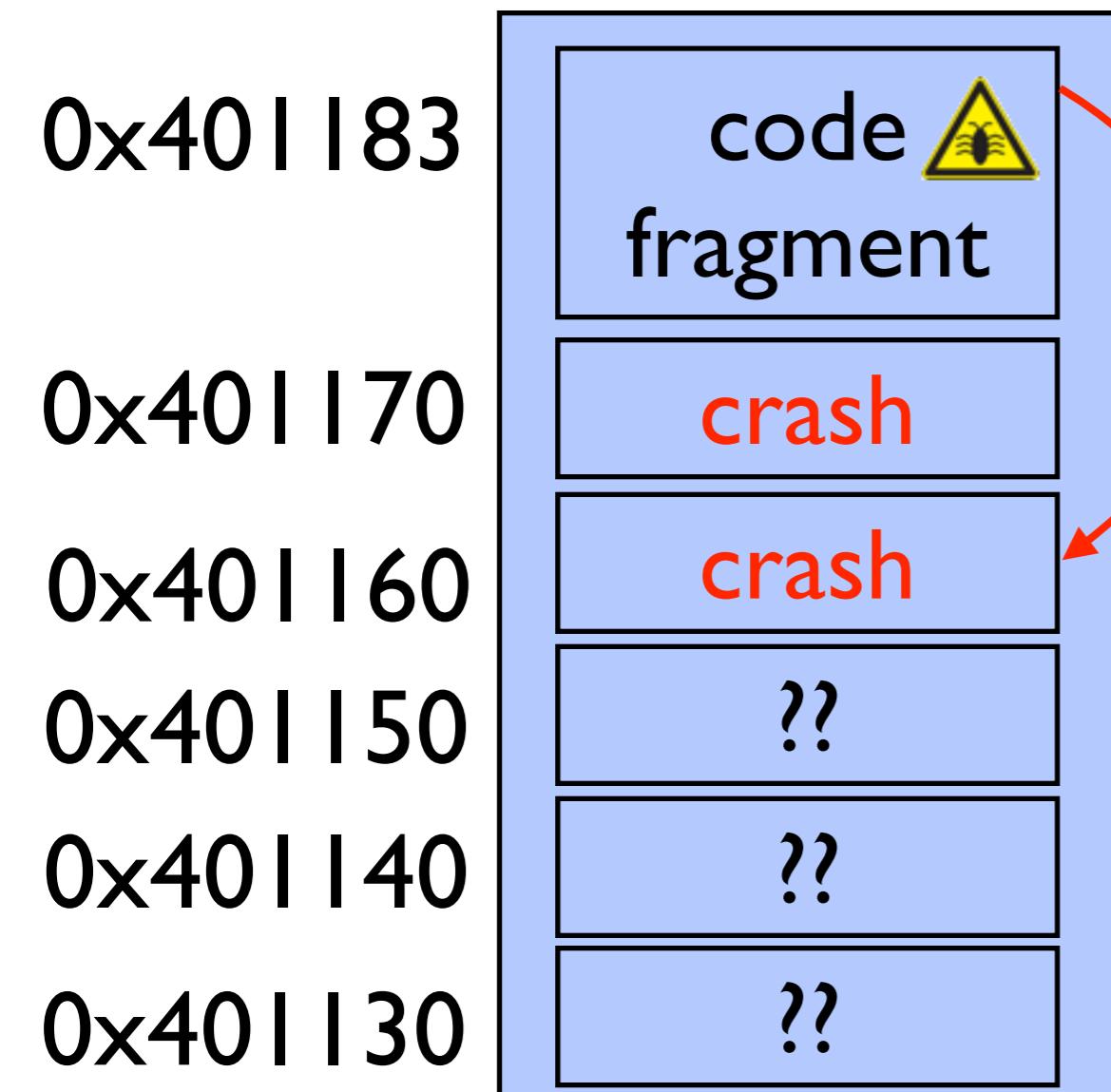
Stack:

return address
0x401170

AAAAAAAAAA
AAAAAAAAAA

How to find gadgets?

.text:



→ Connection closes

Stack:

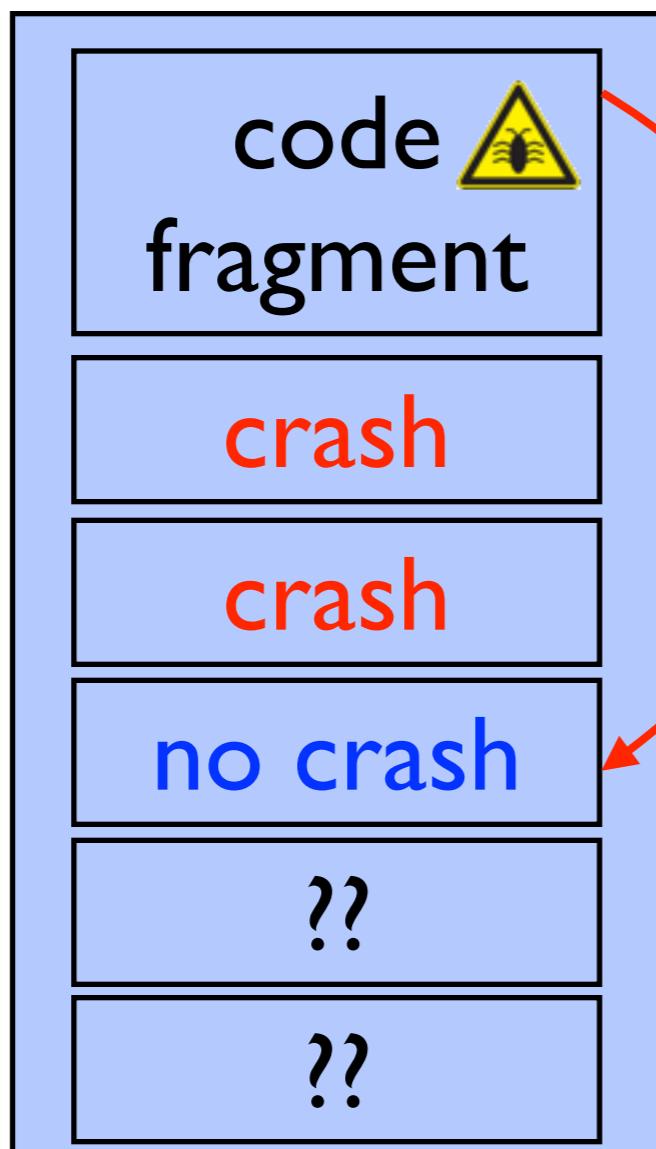
return address
0x401160

AAAAAAAAAA
AAAAAAAAAA

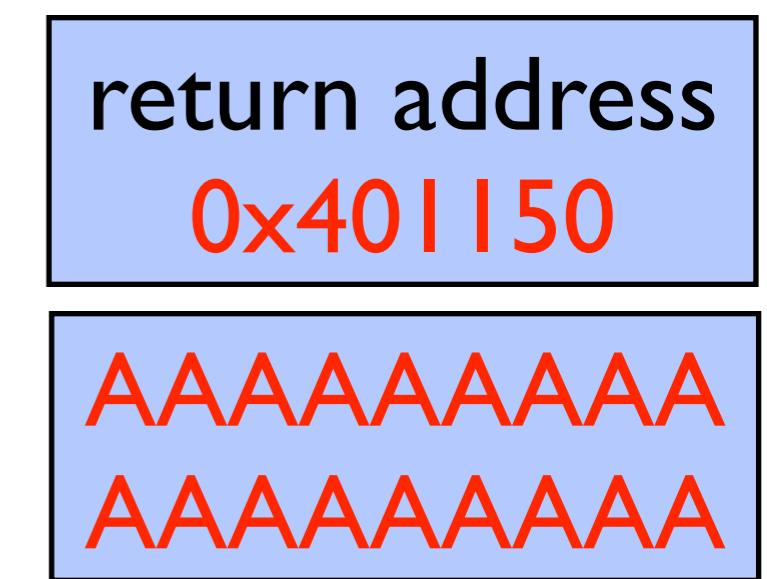
How to find gadgets?

.text:

0x401183
0x401170
0x401160
0x401150
0x401140
0x401130



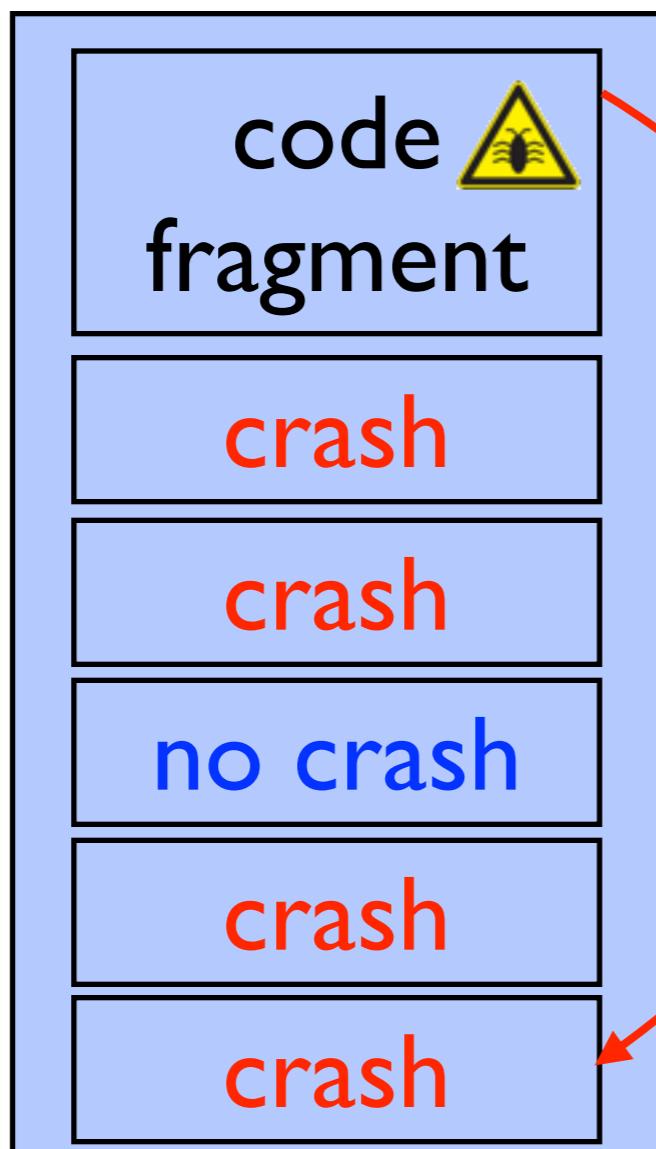
Stack:



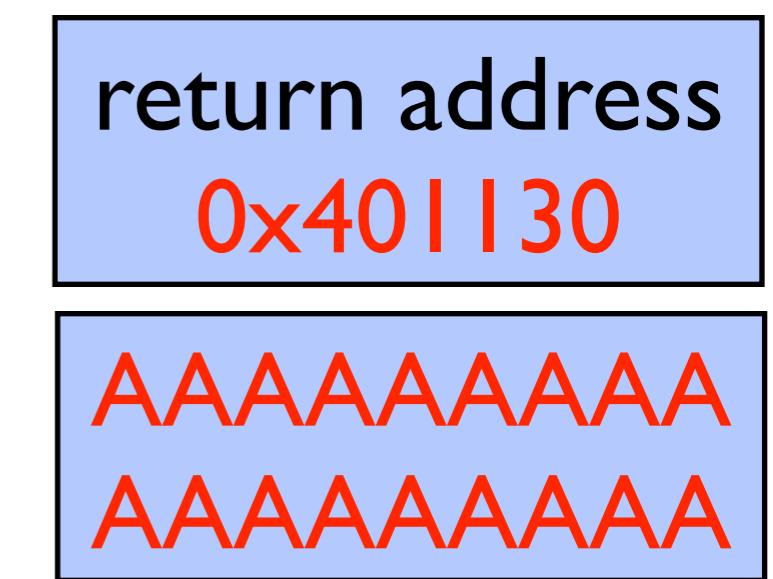
How to find gadgets?

.text:

0x401183
0x401170
0x401160
0x401150
0x401140
0x401130



Stack:



Three types of gadgets

Stop gadget

```
sleep(10);  
return;
```

Crash gadget

```
abort();  
return;
```

Useful gadget

```
dup2(sock, 0);  
return;
```

- Never crashes
- Always crashes
- Crash depends on return

Three types of gadgets

Stop gadget

```
sleep(10);  
return;
```

Crash gadget

```
abort();  
return;
```

Useful gadget

```
dup2(sock, 0);  
return;
```

- Never crashes
- Always crashes
- Crash depends on return

Finding useful gadgets

0x401170

```
dup2(sock, 0);  
return;
```

0x401150

```
sleep(10);  
return;
```

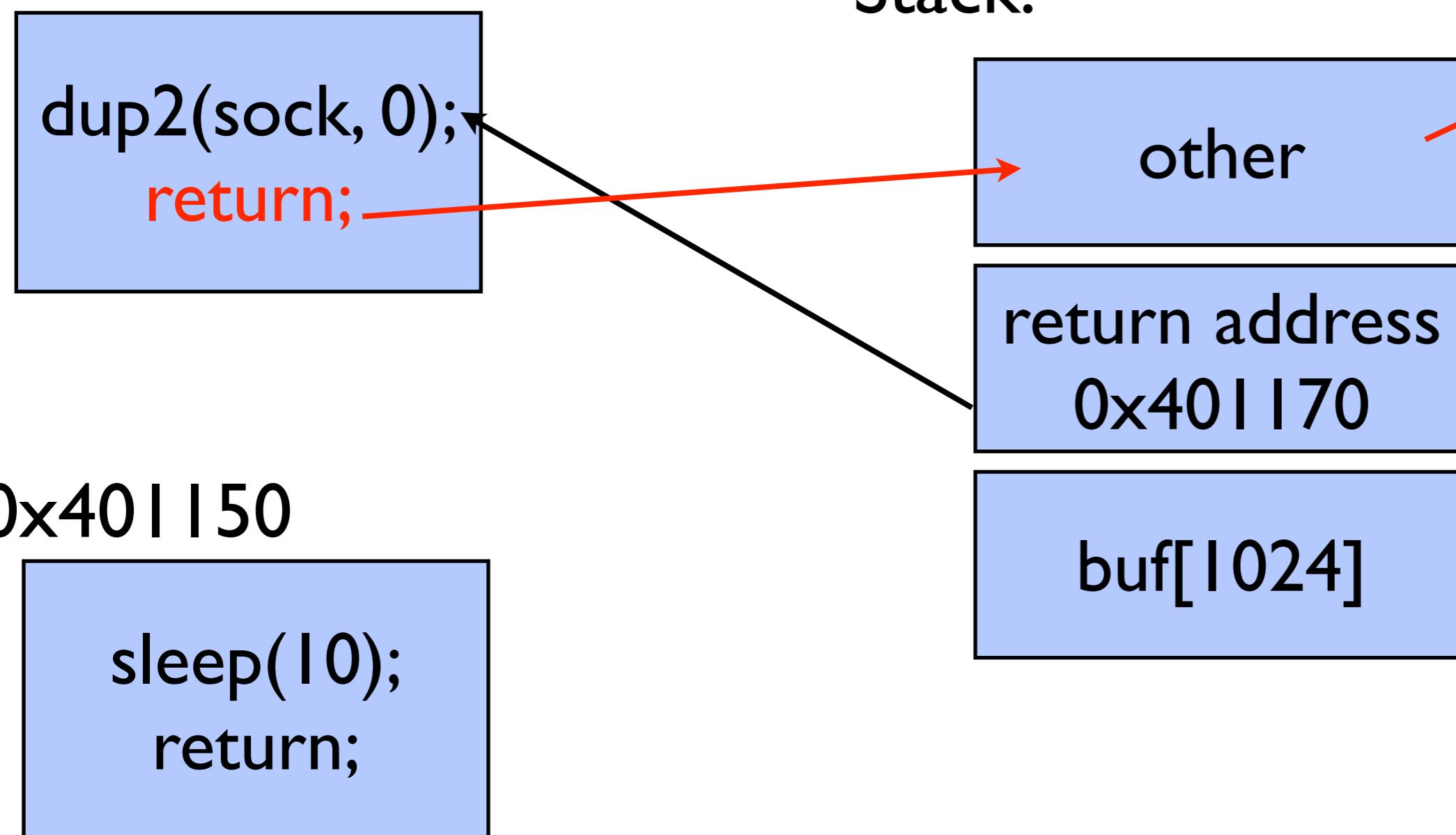
Stack:

other

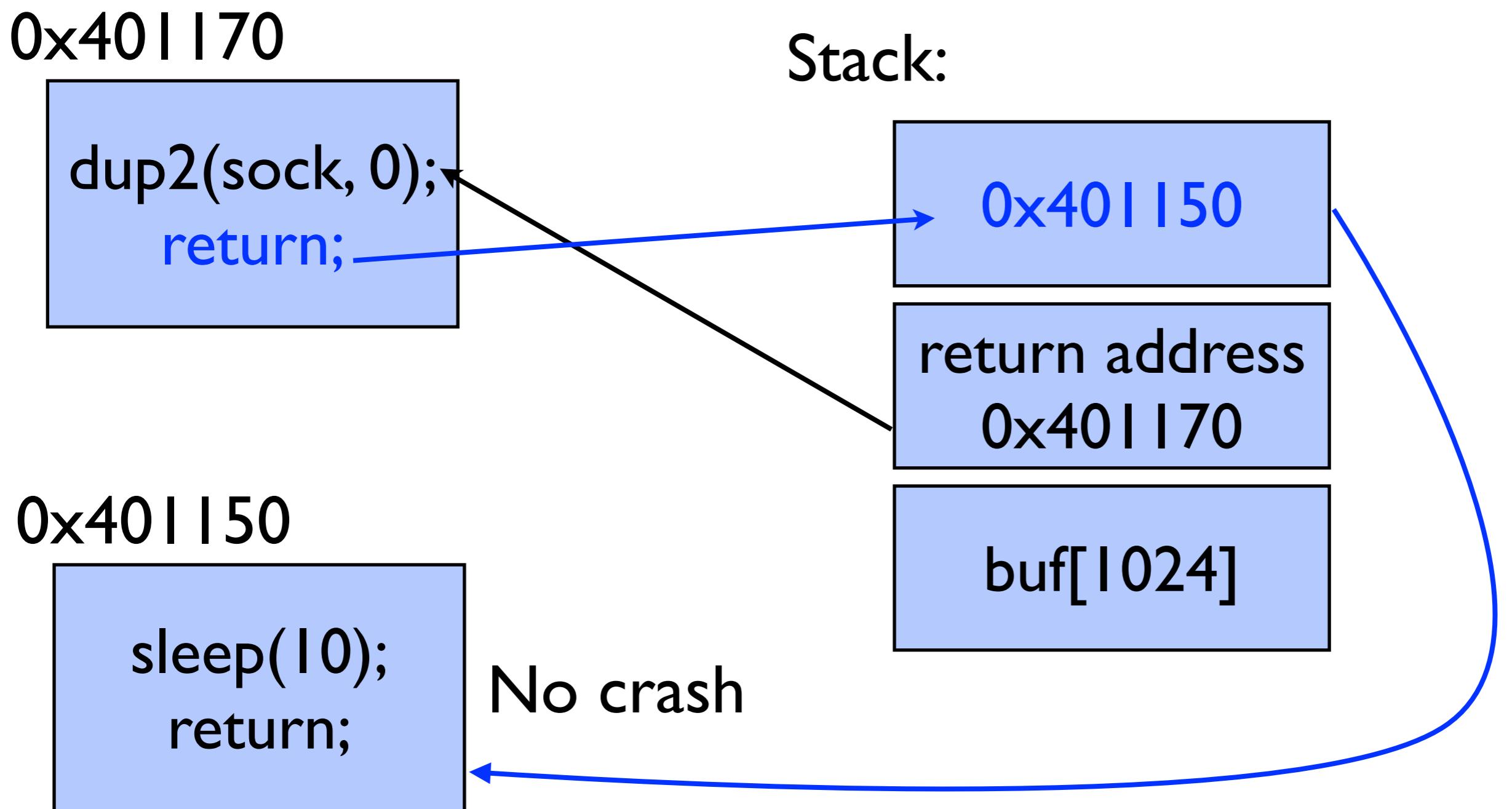
return address
0x401170

buf[1024]

Crash!!

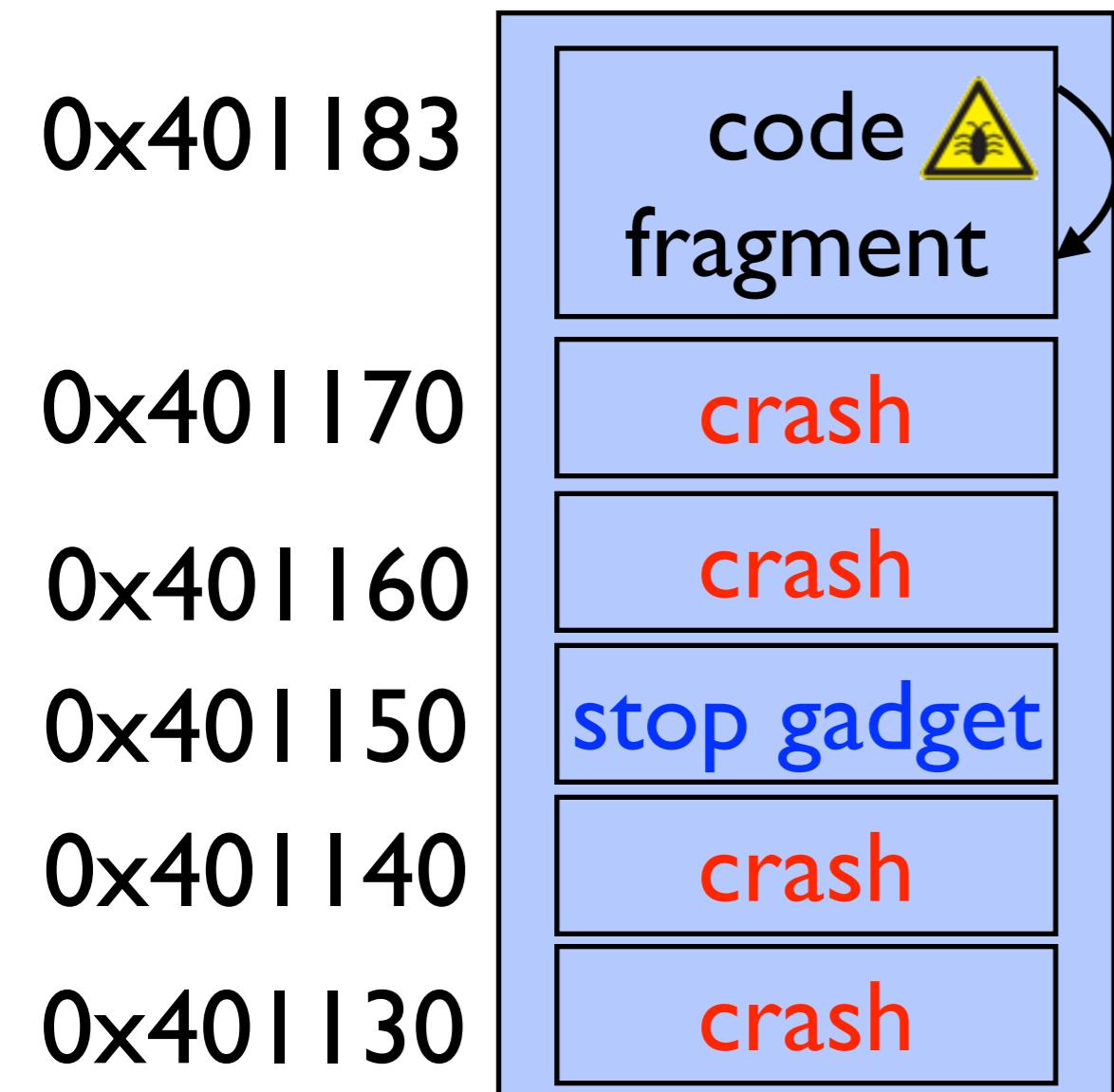


Finding useful gadgets

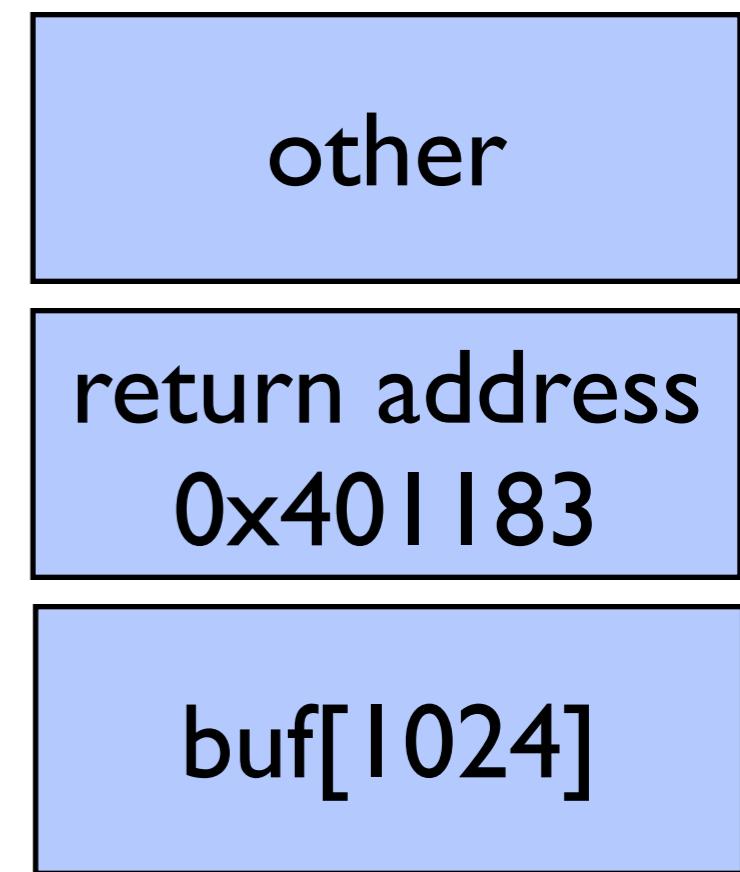


How to find gadgets?

.text:

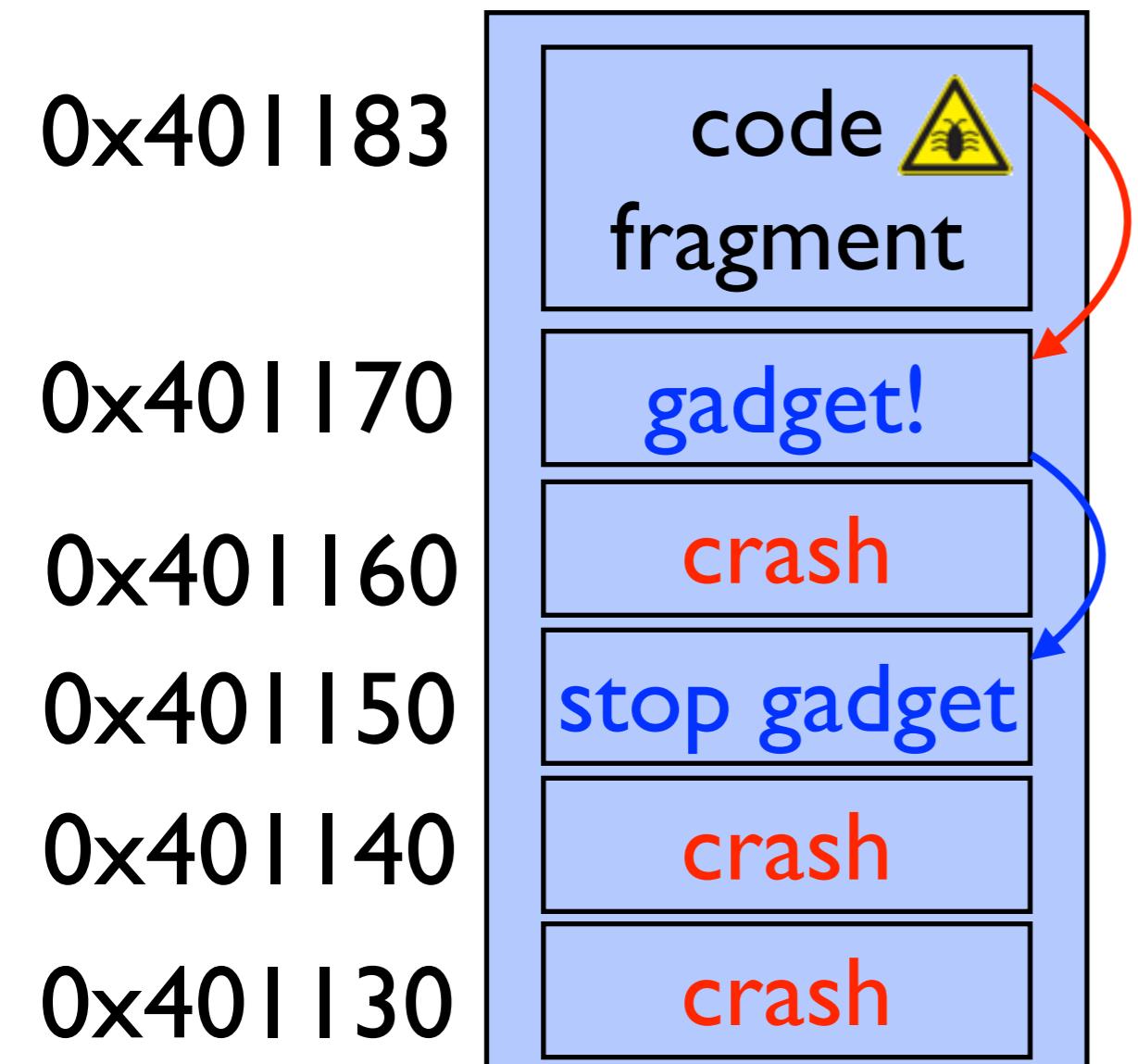


Stack:



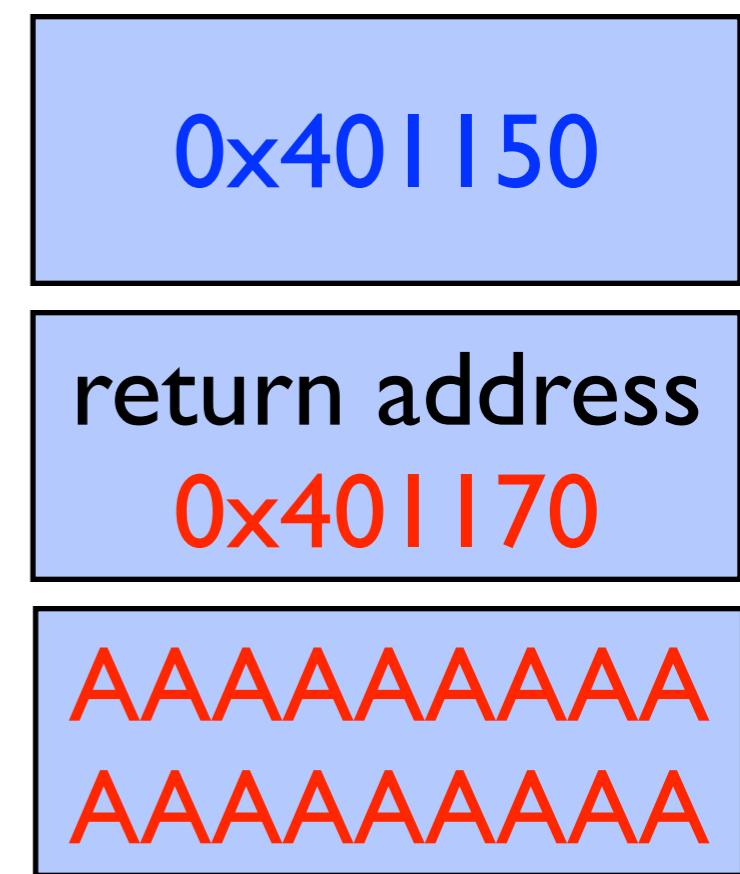
How to find gadgets?

.text:



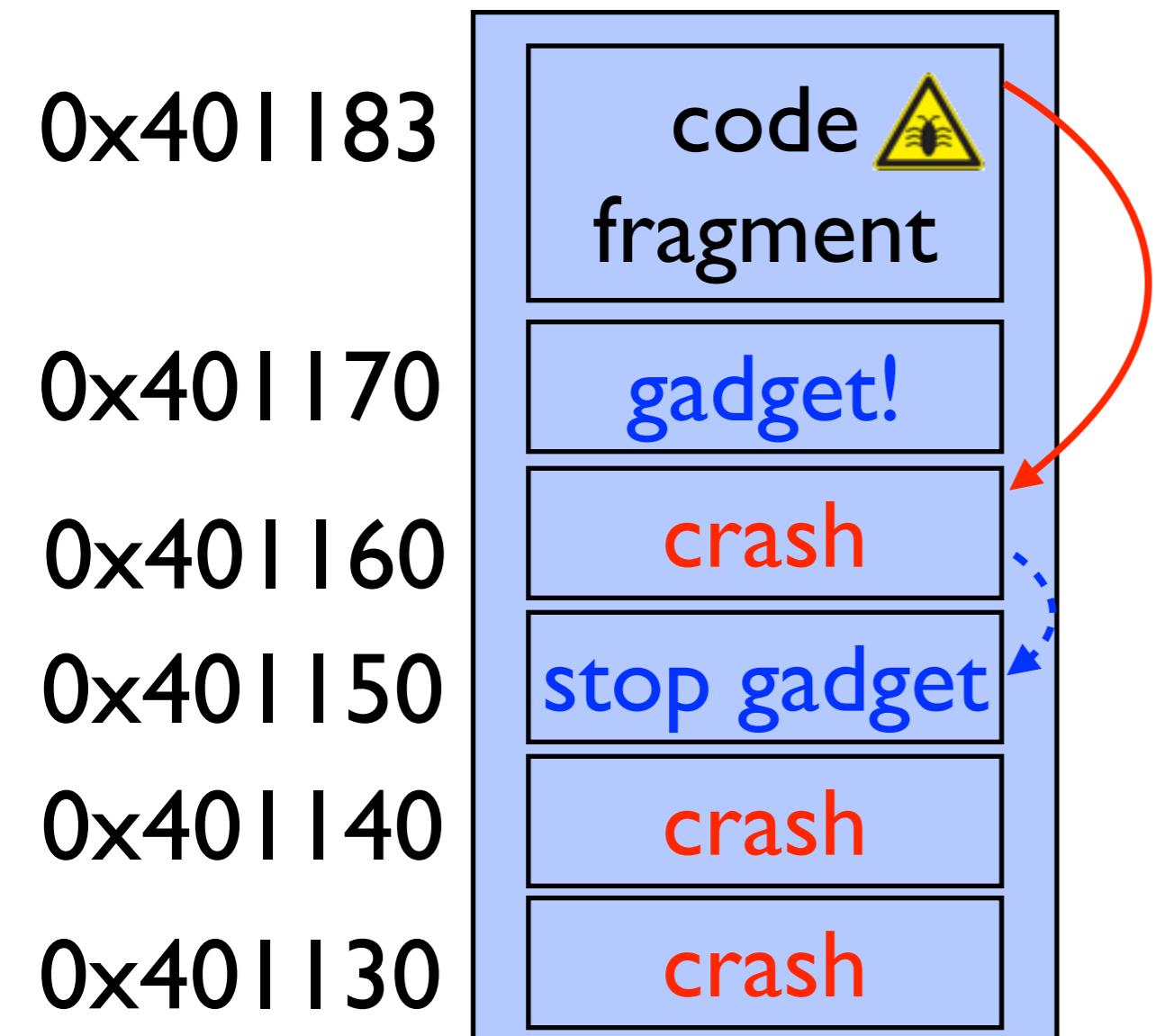
→ Connection hangs

Stack:



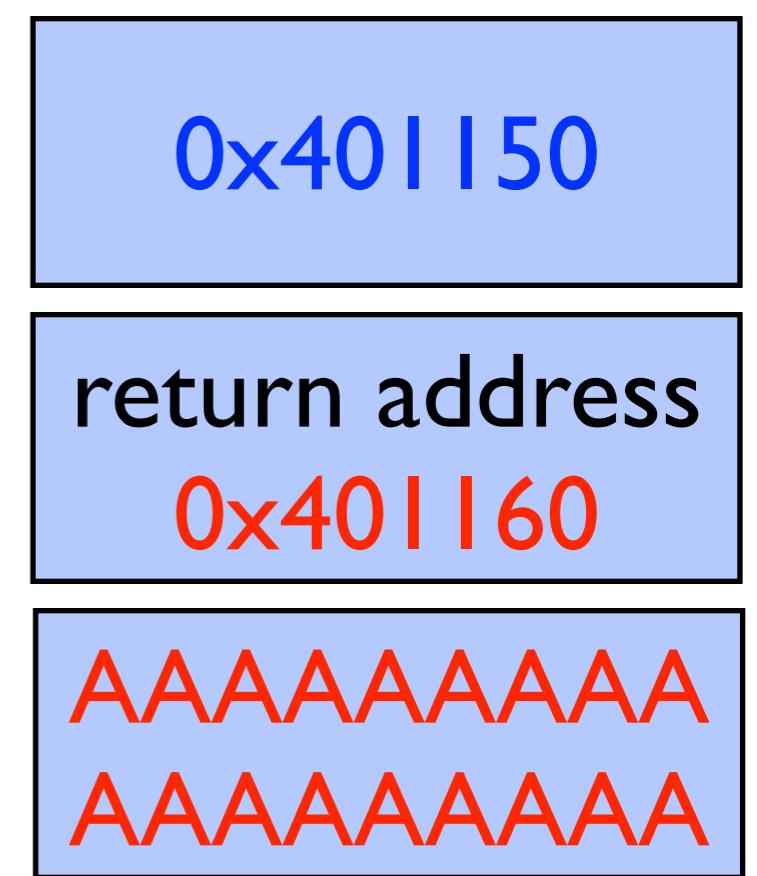
How to find gadgets?

.text:

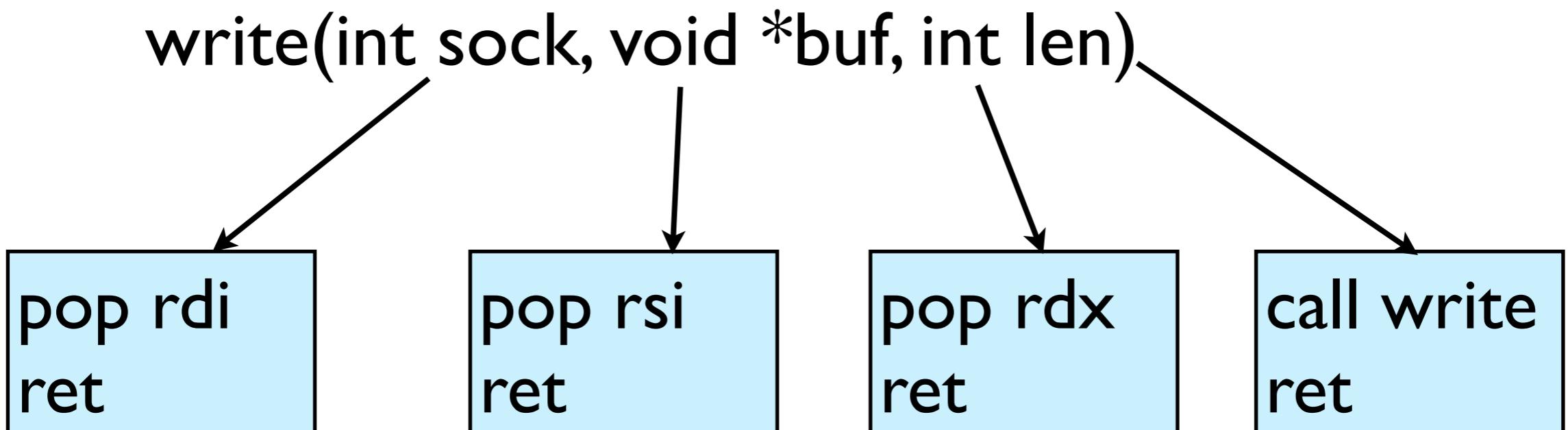


Connection closes

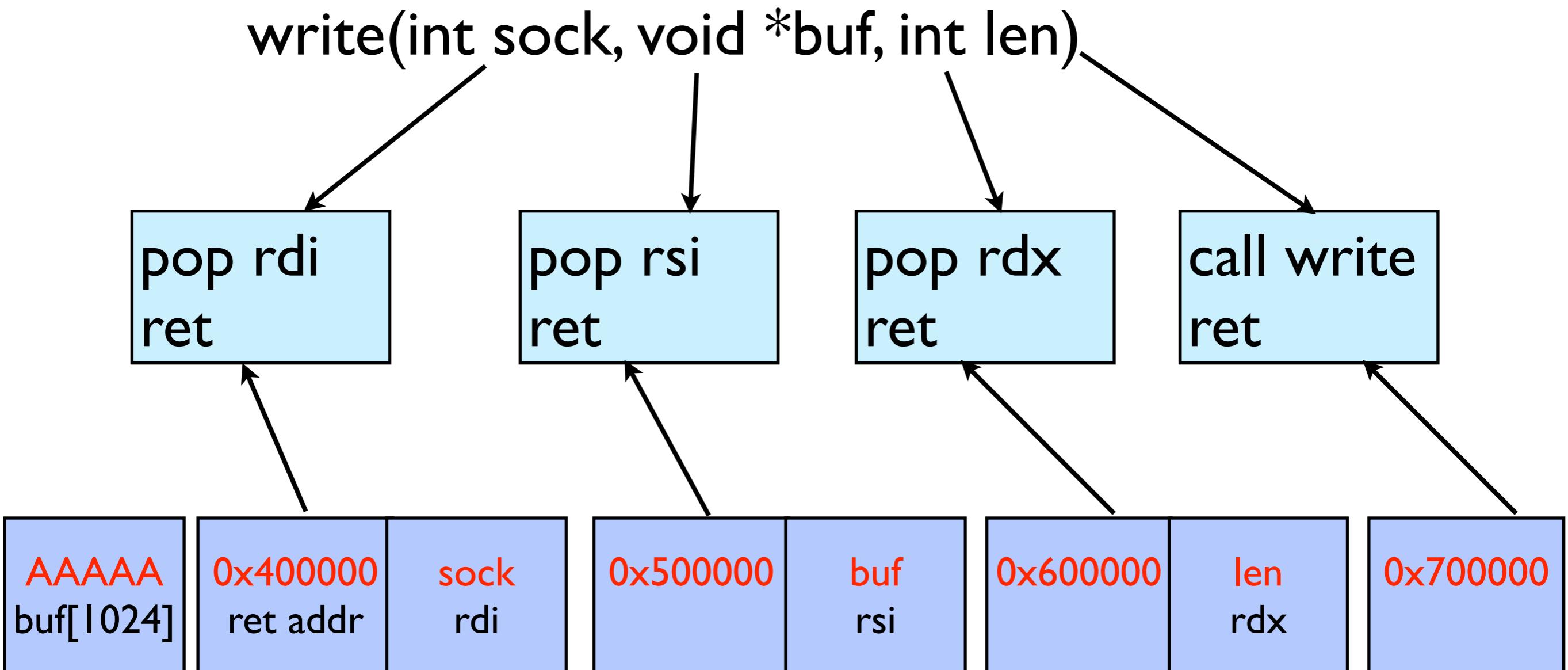
Stack:



What are we looking for?



What are we looking for?



Pieces of the puzzle

pop rsi
ret

pop rdi
ret

pop rdx
ret

call write
ret

stop gadget
[call sleep]

Pieces of the puzzle

The BROP gadget

```
pop rbx  
pop rbp  
pop r12  
pop r13  
pop r14  
pop r15  
ret
```

```
pop rsi  
pop r15  
ret
```

```
pop rdi  
ret
```

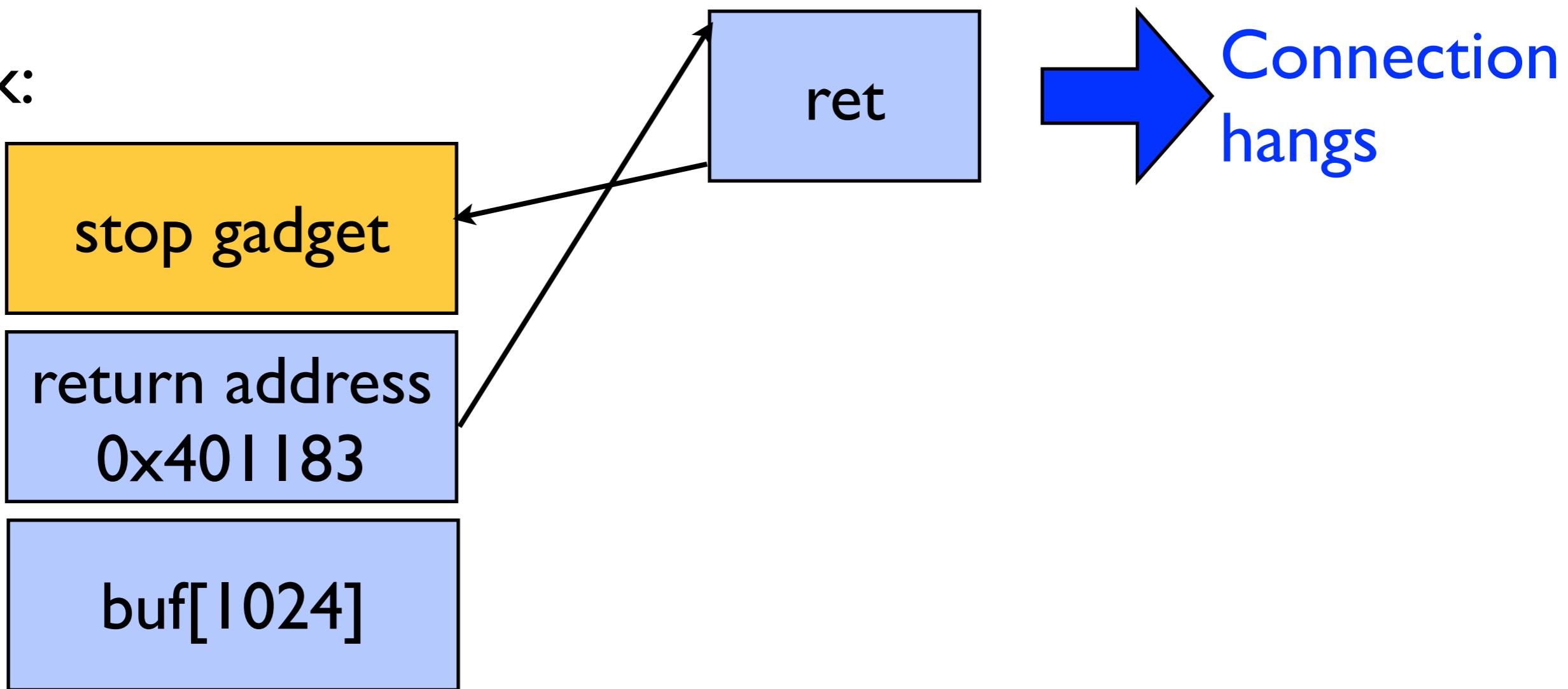
```
pop rdx  
ret
```

```
call write  
ret
```

stop gadget
[call sleep]

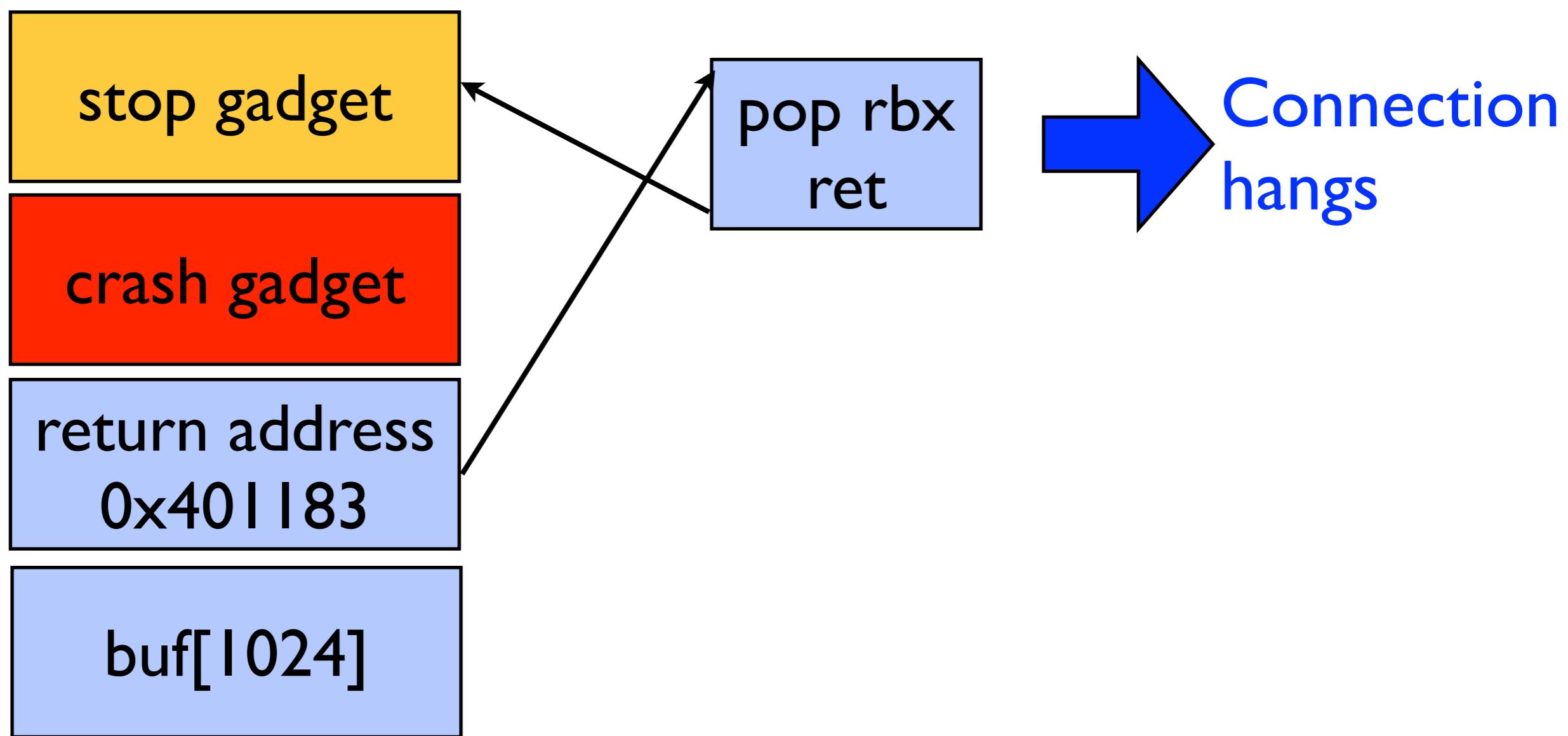
Finding the BROP gadget

Stack:



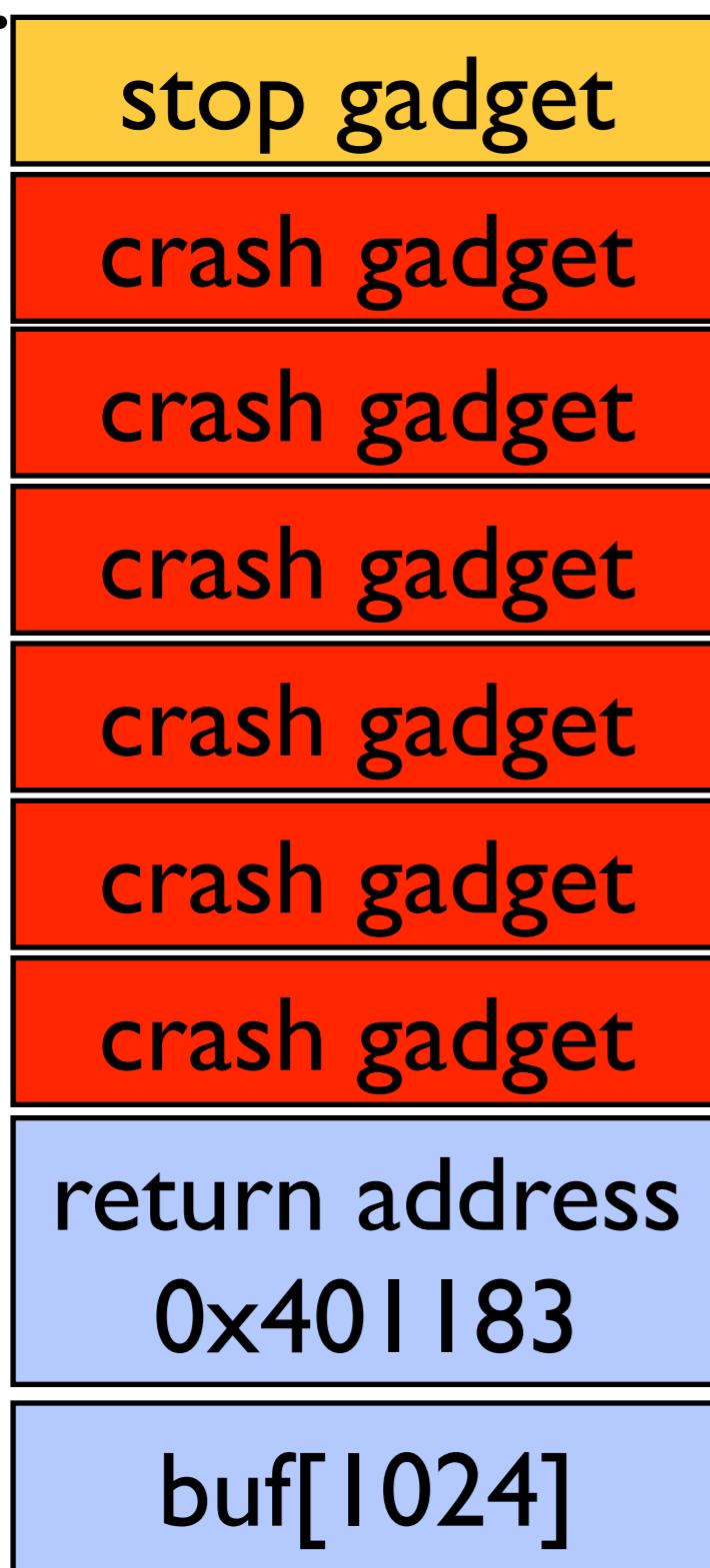
Finding the BROP gadget

Stack:



Finding the BROP gadget

Stack:



pop rbx
pop rbp
pop r12
pop r13
pop r14
pop r15
ret

BROP gadget

Connection
hangs

Pieces of the puzzle

The BROP gadget

```
pop rbx  
pop rbp  
pop r12  
pop r13  
pop r14  
pop r15  
ret
```

```
pop rsi  
pop r15  
ret
```

```
pop rdi  
ret
```

```
pop rdx  
ret
```

```
call write  
ret
```

stop gadget
[call sleep]

Pieces of the puzzle

The BROP gadget

```
pop rbx  
pop rbp  
pop r12  
pop r13  
pop r14  
pop r15  
ret
```

```
pop rsi  
pop r15  
ret
```

```
pop rdi  
ret
```

```
call strcmp  
ret
```

```
call write  
ret
```

stop gadget
[call sleep]

Pieces of the puzzle

The BROP gadget

```
pop rbx  
pop rbp  
pop r12  
pop r13  
pop r14  
pop r15  
ret
```

```
pop rsi  
pop r15  
ret
```

```
pop rdi  
ret
```

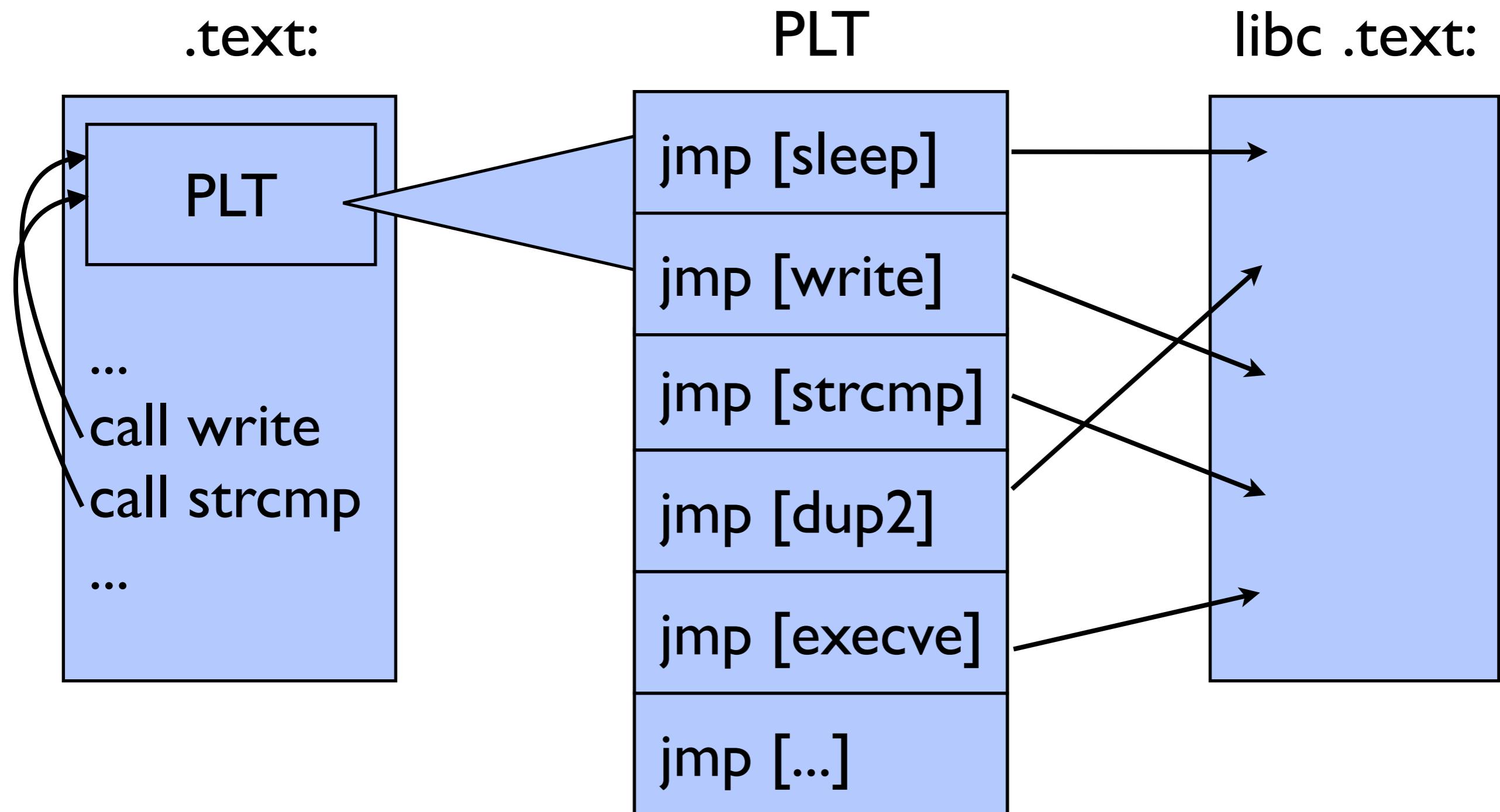
PLT

```
stop gadget  
[call sleep]
```

```
call strcmp  
ret
```

```
call write  
ret
```

Procedure Linking Table (PLT)



Fingerprinting strcmp

arg1	arg2	result
readable	0x0	crash
0x0	readable	crash
readable	readable	nocrash

Can now control three arguments:
strcmp sets RDX to length of string

Finding write

- Try sending data to socket by calling candidate PLT function.
- check if data received on socket.
- chain writes with different FD numbers to find socket. Use multiple connections.

Launching a shell

1. dump binary from memory to network.
Not blind anymore!
2. dump symbol table to find PLT calls.
3. redirect stdin/out to socket:
 - dup2(sock, 0); dup2(sock, 1);
4. read() “/bin/sh” from socket to memory
5. execve(“/bin/sh”, 0, 0)

Braille

- Fully automated: from first crash to shell.
- 2,000 lines of Ruby.
- Needs function that will trigger overflow:
 - nginx: 68 lines.
 - MySQL: 121 lines.
 - toy proprietary service: 35 lines.

try_exp(data) → true crash
false no crash

Attack complexity

