

## Project 2: Thread dispatcher

- **Threads are the most common abstraction for concurrency**
- **Can implement threads inside the OS (“kernel threads”)**
  - Thread create, exit, yield, etc., could be system calls
- **Or can implement threads on top of OS (“Green threads”)**
  - **This is what you will do for this project**
  - Simplifying matters, you will assume only one CPU core
- **OS kernels themselves typically contain threads**
  - Each process managed by a corresponding thread in the kernel
  - Implementation of intra-kernel threads much like Green threads
  - So project will also explain part of how kernels work
- **Goal of project: Implement simple preemptive FIFO scheduler**

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## Background: C++ lambda expressions

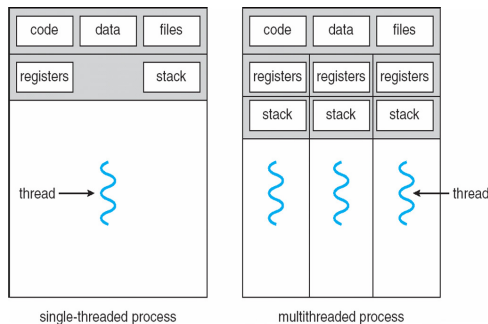
`[captures] (args) {body}` (inferred return type)  
`[captures] (args) -> return-type {body}`

- **Lambda expressions create an unnamed function-like object**
  - Example: `[] (int a, int b) -> int { return a+b; }`
  - Or implicit return type: `[] (int a, int b) { return a+b; }`
- **Can also capture enclosing variables and refer to them**

```
int x = 7, y = 1;
// Capture x by reference, capture y by const copy
auto fn = [&x,y]() { x += y; };
y = 1000;
fn(); // x is now 8, because y was copied into lambda
```
- **Can cast lambda to `std::function<ret(args...)>` type**
  - `std::function<int(int,int)> adder = [] (int a,int b){ return a+b; };`
- **Demo: lambda.cc**

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## Threads vs. processes

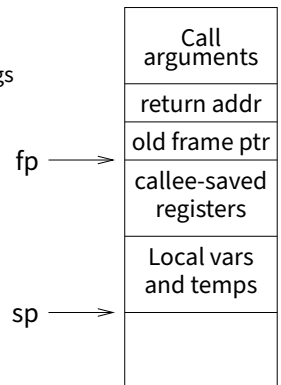


- **A thread is a schedulable execution context**
  - Program counter, stack, registers, ...
- **A process contains an address space, file descriptors, plus one or more threads executing instructions**
- **Demo: thread-v-process.cc**

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## Background: Calling conventions

- **Registers divided into 2 groups**
  - Functions free to clobber *caller-saved* regs
  - But must restore *callee-saved* ones to original value upon return
- **sp register always base of stack**
  - Frame pointer (*fp*) is old *sp* (unless optimized out)
- **Local variables stored in registers and on stack**
- **Function arguments go in caller-saved regs and on stack**



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## Procedure calls

### Procedure call

```
save active caller registers
place arguments in registers/stack
call foo (pushes pc)
    save needed callee registers
    ...do stuff...
    restore callee saved registers
    jump back to calling function
restore stack+caller regs.
```

- **Caller must save some state across function call**
  - Return address, caller-saved registers
- **Other state does not need to be saved**
  - Callee-saved regs, global variables, stack pointer
- **Demo: procedure.cc**

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## Threads vs. procedures

- **Threads may resume out of order**
  - Cannot use LIFO stack to save state
- **General solution: one stack per thread**
  - Switching threads amounts to switching stacks

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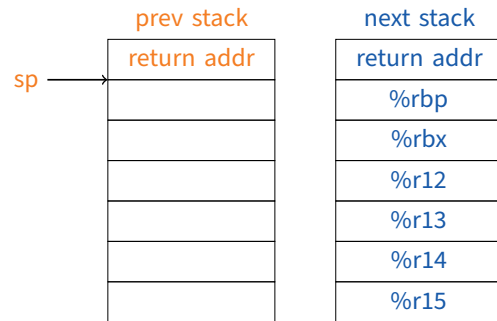
## Stack switching

- Each thread has its own call stack pointed to by `sp`
- How to implement stack switch?
  - Push all registers to save onto stack
  - Save the stack pointer where you can find it again
  - Set the stack pointer to saved value from another stack
  - Pop all saved registers from stack (except now from next stack)
  - Return (except to return address from next stack)
- How to create a new stack?
  - Allocate memory for it (nothing special about memory required)
  - Synthesize a call stack that returns to startup function
  - Make sure startup destroys thread if main function returns
- Project gives you these functions:

```
using sp_t = std::uintptr_t *;
sp_t stack_init(void *stack, size_t bytes, void(*start)());
void stack_switch(sp_t *prev_sp, const sp_t *next_sp);
```

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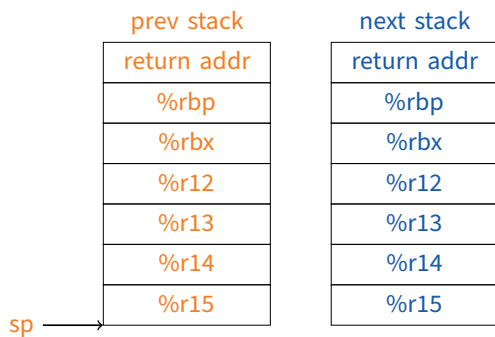
## stack\_switch function



- Demo: `coroutine.cc`

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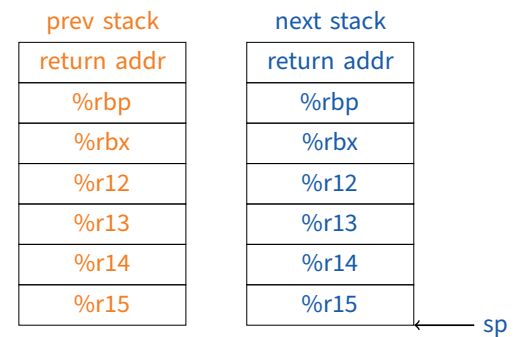
## stack\_switch function



- Demo: `coroutine.cc`

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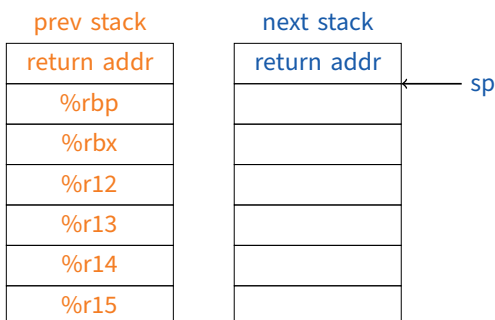
## stack\_switch function



- Demo: `coroutine.cc`

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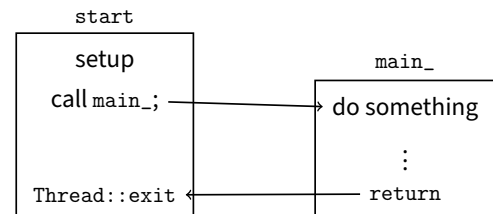
## stack\_switch function



- Demo: `coroutine.cc`

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## Thread initialization



- `stack_init` takes a `void(*)()` (start function)
  - Start function part of your library, same for all threads
  - Must set up environment and deal with thread returning
  - No lambda, so no captures; global current thread may be useful
- `Thread::create` takes a `std::function<void()>` (thread main)
  - Thread main supplied by user of thread library
  - Allows users to initialize threads with lambdas

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## Cooperative threads

- Thread **class bundles all info about thread**
  - Mandatory: stack pointer
  - Possibly useful: memory for stack (for deleting), initial `std::function` to run, pointers for queues
- Global variable with current thread is useful to have
- At lowest level, to switch threads switch stacks
- Demo: `threadlet.cc`
- To create, delete, block, unblock threads, will want queue
  - Can use `std::queue<Thread*>` or keep global head/tail pointers
- Don't delete the memory of the current thread's stack
  - Okay to keep one garbage stack around you delete next time

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## Preemption

```
void timer_init(uint64_t usec, std::function<void()> handler);
bool intr_enabled();
void intr_enable(bool on);
class IntrGuard;
```

- `timer_init` causes timer handler to be called periodically
- Warning: beware of data races!
  - Disable interrupts when touching data used by multiple threads
  - Global data includes `curthread` and scheduler queue
- `IntrGuard` makes it easy to disable interrupts
  - Merely creating an `IntrGuard` object disables interrupts
  - Destroying the `IntrGuard` restores interrupts to previous state
- Demo: `preempt.cc`

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