

Administrivia

- Lab 1 due Friday 10am (5pm if you attend section)
- We will give short extensions to groups that run into trouble. But email us:
 - How much is done and left?
 - How much longer do you need?
- Attend section Friday at 10am to learn about lab 2

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Virtual memory

- Came out of work in late 1960s by **Peter Denning** (lower right)
 - Established working set model
 - Led directly to virtual memory



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Want processes to co-exist

OS	0x9000
gcc	0x7000
bochs/pintos	0x4000
emacs	0x3000
	0x0000

- Consider multiprogramming on physical memory
 - What happens if pintos needs to expand?
 - If emacs needs more memory than is on the machine?
 - If pintos has an error and writes to address 0x7100?
 - When does gcc have to know it will run at 0x4000?
 - What if emacs isn't using its memory?

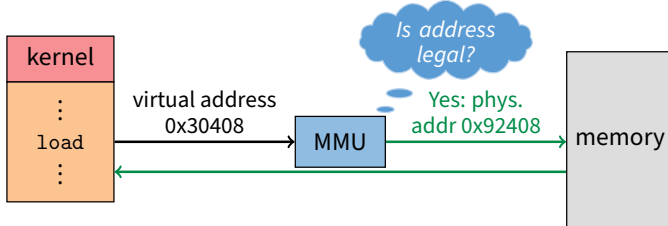
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Issues in sharing physical memory

- **Protection**
 - A bug in one process can corrupt memory in another
 - Must somehow prevent process A from trashing B's memory
 - Also prevent A from even observing B's memory (ssh-agent)
- **Transparency**
 - A process shouldn't require particular physical memory bits
 - Yet processes often require large amounts of contiguous memory (for stack, large data structures, etc.)
- **Resource exhaustion**
 - Programmers typically assume machine has "enough" memory
 - Sum of sizes of all processes often greater than physical memory

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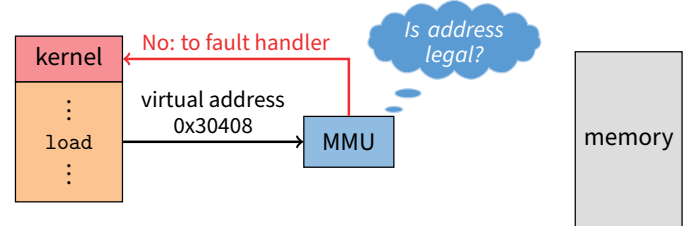
Virtual memory goals



- Give each program its own *virtual* address space
 - At runtime, *Memory-Management Unit* relocates each load/store
 - Application doesn't see *physical* memory addresses
- Also enforce protection
 - Prevent one app from messing with another's memory
- And allow programs to see more memory than exists
 - Somehow relocate some memory accesses to disk

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Virtual memory goals

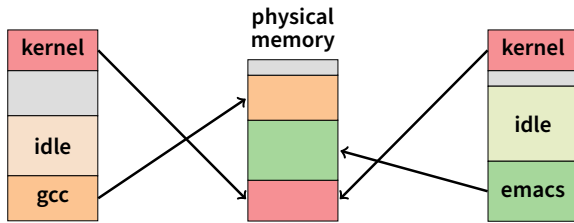


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Virtual memory advantages

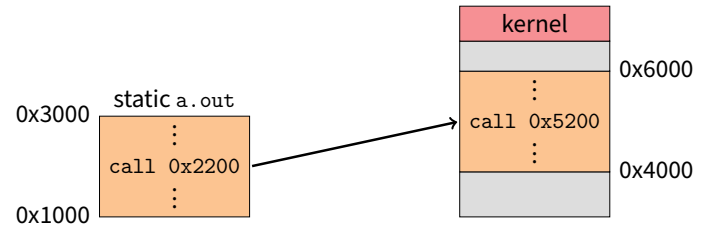
- Can re-locate program while running
 - Run partially in memory, partially on disk
- Most of a process's memory may be idle (80/20 rule).



- Write idle parts to disk until needed
- Let other processes use memory of idle part
- Like CPU virtualization: when process not using CPU, switch (Not using a memory region? switch it to another process)
- Challenge: VM = extra layer, could be slow

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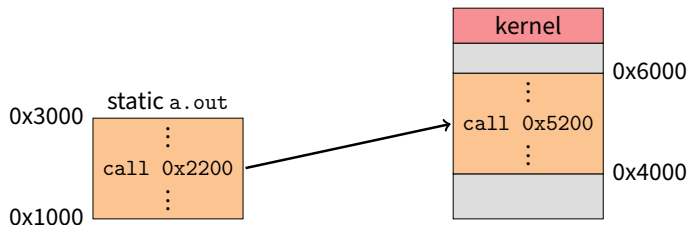
Idea 1: load-time linking



- **Linker** patches addresses of symbols like `printf`
- **Idea: link when process executed, not at compile time**
 - Determine where process will reside in memory
 - Adjust all references within program (using addition)
- **Problems?**

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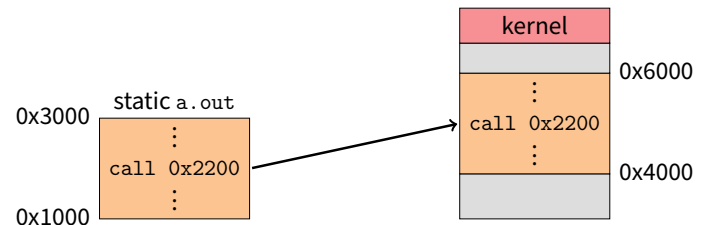
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- **Linker** patches addresses of symbols like `printf`
- **Idea: link when process executed, not at compile time**
 - Determine where process will reside in memory
 - Adjust all references within program (using addition)
- **Problems?**
 - How to enforce protection?
 - How to move once already in memory? (consider data pointers)
 - What if no contiguous free region fits program?

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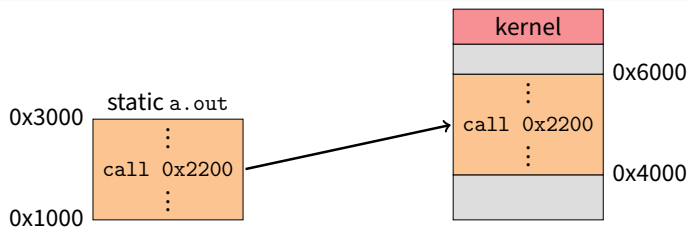
Idea 2: base + bound register



- **Two special privileged registers: base and bound**
- **On each load/store/jump:**
 - Physical address = virtual address + **base**
 - Check $0 \leq \text{virtual address} < \text{bound}$, else trap to kernel
- **How to move process in memory?**
- **What happens on context switch?**

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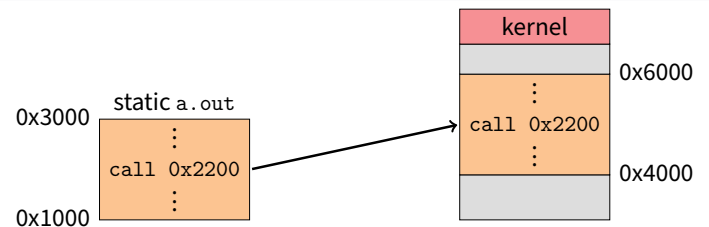
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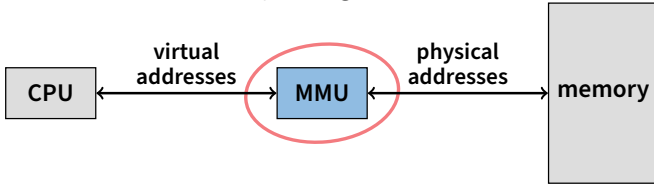


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- **How to move process in memory?**
 - Change **base** register
- **What happens on context switch?**
 - OS must re-load **base** and **bound** register

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Definitions

- Programs load/store to **virtual addresses**
- Actual memory uses **physical addresses**
- VM Hardware is Memory Management Unit (**MMU**)

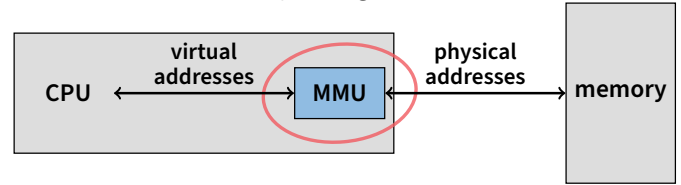


- Usually part of CPU
- Configured through privileged instructions (e.g., load bound reg)
- Translates from virtual to physical addresses
- Gives per-process view of memory called **address space**

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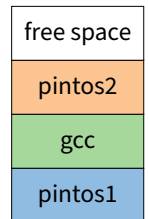
Base+bound trade-offs

- **Advantages**
 - Cheap in terms of hardware: only two registers
 - Cheap in terms of cycles: do add and compare in parallel
 - Examples: Cray-1 used this scheme
- **Disadvantages**

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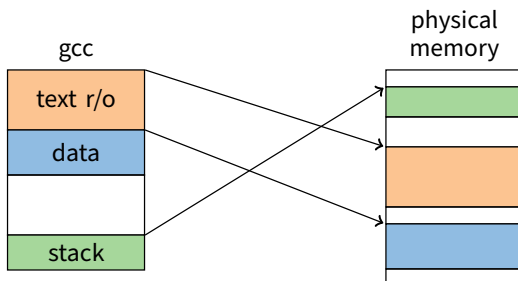
Base+bound trade-offs

- **Advantages**
 - Cheap in terms of hardware: only two registers
 - Cheap in terms of cycles: do add and compare in parallel
 - Examples: Cray-1 used this scheme
- **Disadvantages**
 - Growing a process is expensive or impossible
 - No way to share code or data (E.g., two copies of bochs, both running pintos)
- **One solution: Multiple segments**
 - E.g., separate code, stack, data segments
 - Possibly multiple data segments



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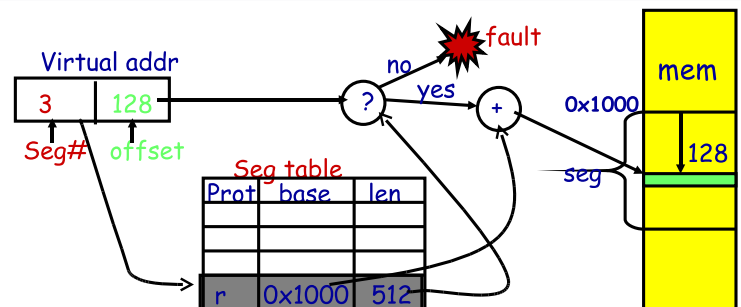
Segmentation



- **Let processes have many base/bound regs**
 - Address space built from many segments
 - Can share/protect memory at segment granularity
- **Must specify segment as part of virtual address**

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Segmentation mechanics

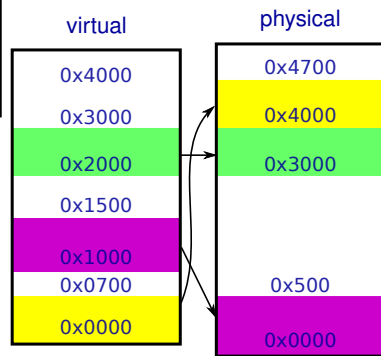


- **Each process has a segment table**
- **Each VA indicates a segment and offset:**
 - Top bits of addr select segment, low bits select offset (PDP-10)
 - Or segment selected by instruction or operand (means you need wider "far" pointers to specify segment)

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Segmentation example

Seg	base	bounds	rw
0	0x4000	0x6fff	10
1	0x0000	0x4fff	11
2	0x3000	0xffff	11
3			00



- **2-bit segment number (1st digit), 12 bit offset (last 3)**
 - Where is 0x0240? 0x1108? 0x265c? 0x3002? 0x1600?

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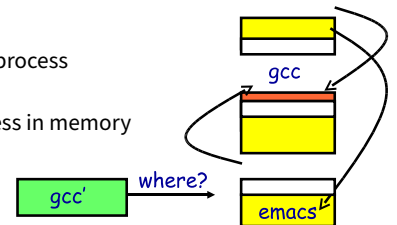
Segmentation trade-offs

Advantages

- Multiple segments per process
- Allows sharing! (how?)
- Don't need entire process in memory

Disadvantages

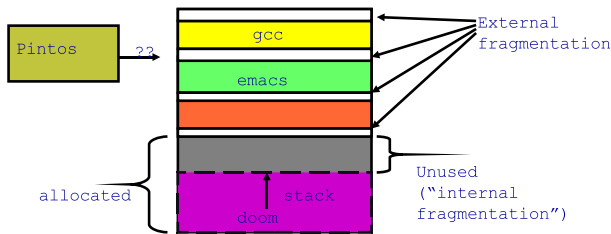
- Requires translation hardware, which could limit performance
- Segments not completely transparent to program (e.g., default segment faster or uses shorter instruction)
- n byte segment needs n contiguous bytes of physical memory
- Makes *fragmentation* a real problem.



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Fragmentation

- **Fragmentation** \implies Inability to use free memory
- **Over time:**
 - Variable-sized pieces = many small holes (external fragmentation)
 - Fixed-sized pieces = no external holes, but force internal waste (internal fragmentation)



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Alternatives to hardware MMU

Language-level protection (JavaScript)

- Single address space for different modules
- Language enforces isolation
- Singularity OS does this with C# [Hunt]

Software fault isolation

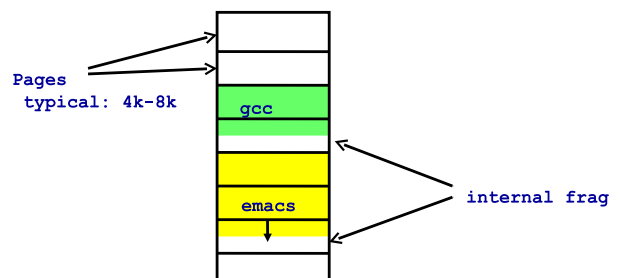
- Instrument compiler output
- Checks before every store operation prevents modules from trashing each other
- Google's now deprecated [Native Client](#) does this for x86 [Yee]
- Easier to do for virtual architecture, e.g., [Wasm](#)

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Paging

- Divide memory up into small, equal-size *pages*
- Map virtual pages to physical pages
 - Each process has separate mapping
- Allow OS to gain control on certain operations
 - Read-only pages trap to OS on write
 - Invalid pages trap to OS on read or write
 - OS can change mapping and resume application
- Other features sometimes found:
 - Hardware can set "accessed" and "dirty" bits
 - Control page execute permission separately from read/write
 - Control caching or memory consistency of page

Paging trade-offs

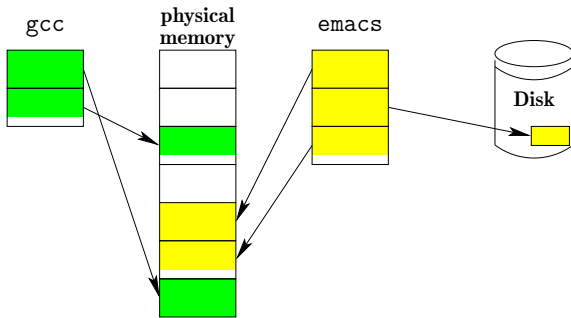


- Eliminates external fragmentation
- Simplifies allocation, free, and backing storage (swap)
- Average internal fragmentation of .5 pages per "segment"

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Simplified allocation

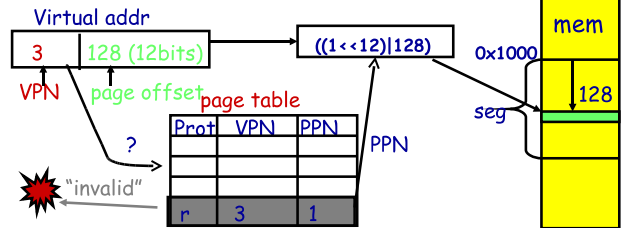


- Allocate any physical page to any process
- Can store idle virtual pages on disk

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Paging data structures

- Pages are fixed size, e.g., 4 KiB
 - Least significant 12 ($\log_2 4 \text{ Ki}$) bits of address are *page offset*
 - Most significant bits are *page number*
- Each process has a *page table*
 - Maps *virtual page numbers (VPNs)* to *physical page numbers (PPNs)*
 - Also includes bits for protection, validity, etc.
- On memory access: Translate VPN to PPN, then add offset



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Example: Paging on PDP-11

- 64 KiB virtual memory, 8 KiB pages
 - Separate address space for instructions & data
 - I.e., can't read your own instructions with a load
- Entire page table stored in registers
 - 8 Instruction page translation registers
 - 8 Data page translations
- Swap 16 machine registers on each context switch

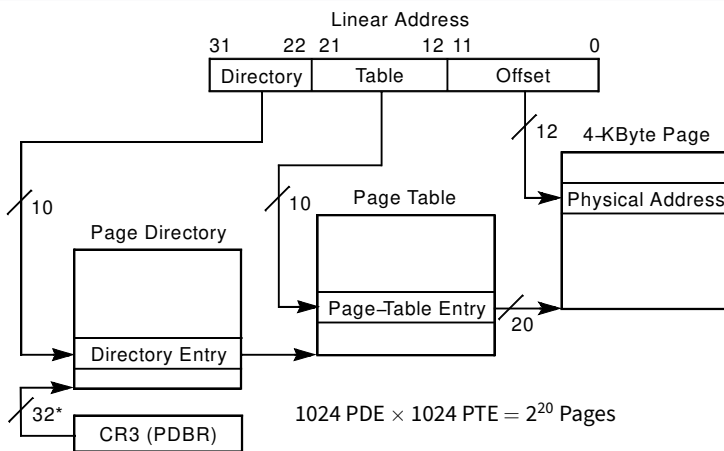
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x86 Paging

- Paging enabled by bits in a control register (`%cr0`)
 - Only privileged OS code can manipulate control registers
- Normally 4 KiB pages
- `%cr3`: points to physical address of 4 KiB page directory
 - See `pagedir_activate` in Pintos
- Page directory: 1024 PDEs (page directory entries)
 - Each contains physical address of a page table
- Page table: 1024 PTEs (page table entries)
 - Each contains physical address of virtual 4K page
 - Page table covers 4 MiB of Virtual mem
- See [old intel manual](#) for simplest explanation
 - Also volume 2 of [AMD64 Architecture docs](#)
 - Also volume 3A of [latest intel 64 architecture manual](#)

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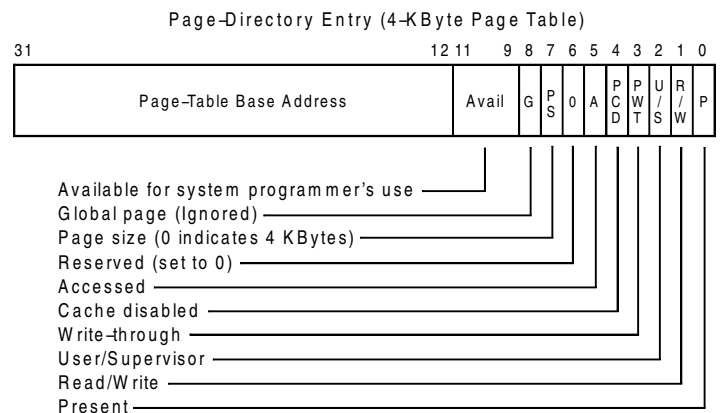
x86 page translation



*32 bits aligned onto a 4-KByte boundary

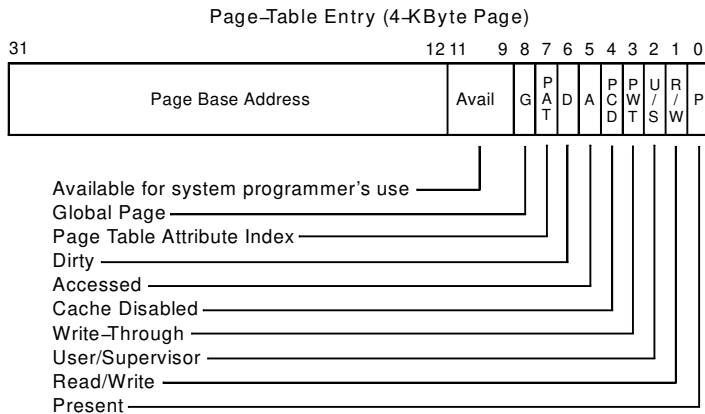
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x86 page directory entry



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x86 page table entry



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x86 hardware segmentation

- x86 architecture *also* supports segmentation
 - Segment register base + pointer val = *linear address*
 - Page translation happens on linear addresses
- Two levels of protection and translation check
 - Segmentation model has four privilege levels (CPL 0–3)
 - Paging only two, so 0–2 = kernel, 3 = user
- Why do you want *both* paging and segmentation?

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 - Segmentation model has four privilege levels (CPL 0–3)
 - Paging only two, so 0–2 = kernel, 3 = user
- Why do you want *both* paging and segmentation?
- Short answer: You don't – just adds overhead
 - Most OSes use "flat mode" – set base = 0, bounds = 0xffffffff in all segment registers, then forget about it
 - x86-64 architecture removes much segmentation support
- Long answer: Has some fringe/incidental uses
 - Keep pointer to thread-local storage w/o wasting normal register
 - 32-bit VMware runs guest OS in CPL 1 to trap stack faults
 - OpenBSD used CS limit for W^X when no PTE NX bit

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Making paging fast

- x86 PTs require 3 memory references per load/store
 - Look up page table address in page directory
 - Look up physical page number (PPN) in page table
 - Actually access physical page corresponding to virtual address
- For speed, CPU caches recently used translations
 - Called a *translation lookaside buffer* or **TLB**
 - Typical: 64-2k entries, 4-way to fully associative, 95% hit rate
 - Modern CPUs add second-level TLB with ~1,024+ entries; often separate instruction and data TLBs
 - Each TLB entry maps a VPN → PPN + protection information
- On each memory reference
 - Check TLB, if entry present get physical address fast
 - If not, walk page tables, insert in TLB for next time (Must evict some entry)

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TLB details

- TLB operates at CPU pipeline speed ⇒ small, fast
- Complication: what to do when switching address space?
 - Flush TLB on context switch (e.g., old x86)
 - Tag each entry with associated process's ID (e.g., MIPS)
- In general, OS must manually keep TLB valid
 - Changing page table in memory won't affect cached TLB entry
- E.g., on x86 must use *invlpg* instruction
 - Invalidates a page translation in TLB
 - Note: very expensive instruction (100–200 cycles)
 - Must execute after changing a possibly used page table entry
 - Otherwise, hardware will miss page table change
- More Complex on a multiprocessor (TLB shutdown)
 - Requires sending an interprocessor interrupt (IPI)
 - Remote processor must execute *invlpg* instruction

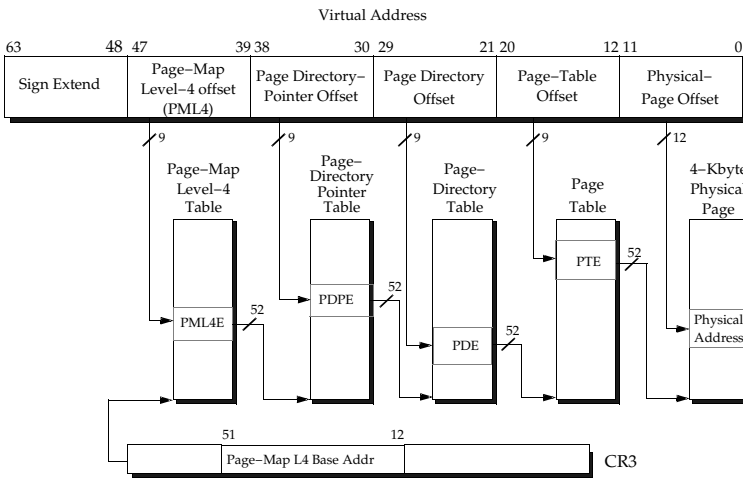
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x86 Paging Extensions

- PSE: Page size extensions
 - Setting bit 7 in PDE makes a 4 MiB translation (no PT)
- PAE Page address extensions
 - Newer 64-bit PTE format allows 36+ bits of physical address
 - Page tables, directories have only 512 entries
 - Use 4-entry Page-Directory-Pointer Table to regain 2 lost bits
 - PDE bit 7 allows 2 MiB translation
- Long mode PAE (x86-64)
 - In Long mode, pointers are 64-bits
 - Extends PAE to map 48 bits of virtual address (next slide)
 - Why aren't all 64 bits of VA usable?

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x86 long mode paging



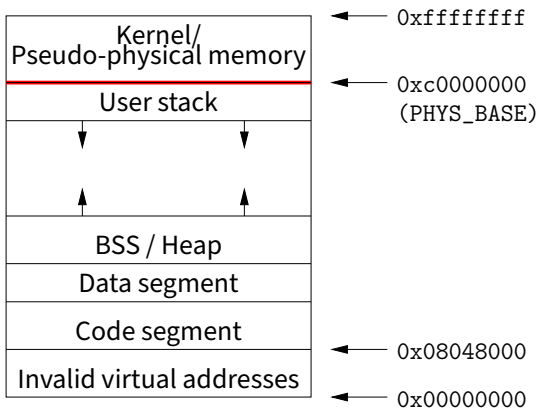
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Where does the OS live?

- **In its own address space?**
 - Can't do this on most hardware (e.g., syscall instruction won't switch address spaces)
 - Also would make it harder to parse syscall arguments passed as pointers
- **So in the same address space as process**
 - Use protection bits to prohibit user code from writing kernel
- **Typically all kernel text, most data at same VA in every address space**
 - On x86, must manually set up page tables for this
 - Usually just map kernel in contiguous virtual memory when boot loader puts kernel into contiguous physical memory
 - Some hardware puts physical memory (kernel-only) somewhere in virtual address space
 - Typically kernel goes in high memory; with signed numbers, can mean small negative addresses (small linker relocations)

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Pintos memory layout



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Very different MMU: MIPS

- **Hardware checks TLB on application load/store**
 - References to addresses not in TLB trap to kernel
- **Each TLB entry has the following fields: Virtual page, Pid, Page frame, NC, D, V, Global**
- **Kernel itself unpagged**
 - All of physical memory contiguously mapped in high VM (hardwired in CPU, not just by convention as with Pintos)
 - Kernel uses these pseudo-physical addresses
- **User TLB fault handler very efficient**
 - Two hardware registers reserved for it
 - utlb miss handler can itself fault—allow paged page tables
- **OS is free to choose page table format!**

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DEC Alpha MMU

- **Firmware managed TLB**
 - Like MIPS, TLB misses handled by software
 - Unlike MIPS, TLB miss routines ship with machine in ROM (but copied to main memory on boot—so can be overwritten)
 - Firmware known as “PAL code” (privileged architecture library)
- **Hardware capabilities**
 - 8 KiB, 64 KiB, 512 KiB, 4 MiB pages all available
 - TLB supports 128 instruction/128 data entries of any size
- **Various other events vector directly to PAL code**
 - call_pal instruction, TLB miss/fault, FP disabled
- **PAL code runs in special privileged processor mode**
 - Interrupts always disabled
 - Have access to special instructions and registers

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PAL code interface details

- **Examples of Digital Unix PALcode entry functions**
 - callsys/retsys - make, return from system call
 - swpctx - change address spaces
 - wrvptptr - write virtual page table pointer
 - tbi - TBL invalidate
- **Some fields in PALcode page table entries**
 - GH - 2-bit granularity hint → 2^N pages have same translation
 - ASM - address space match → mapping applies in all processes

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Example: Paging to disk

- `gcc` needs a new page of memory
- OS re-claims an idle page from `emacs`
- If page is *clean* (i.e., also stored on disk):
 - E.g., page of text from `emacs` binary on disk
 - Can always re-read same page from binary
 - So okay to discard contents now & give page to `gcc`
- If page is *dirty* (meaning memory is only copy)
 - Must write page to disk first before giving to `gcc`
- Either way:
 - Mark page invalid in `emacs`
 - `emacs` will fault on next access to virtual page
 - On fault, OS reads page data back from disk into new page, maps new page into `emacs`, resumes executing

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Paging in day-to-day use

- Demand paging
- Growing the stack
- BSS page allocation
- Shared text
- Shared libraries
- Shared memory
- Copy-on-write (`fork`, `mmap`, etc.)
- Q: Which pages should have global bit set on x86?

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